

Words and Groups

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Additive Number Theory

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- In 1909 Hilbert solved the problem affirmatively.
- **Non-commutative analogues:**
Present group elements as short products of special elements: powers, or commutators, or values of a general word w , or elements of a given conjugacy class in the group.

- Let $w = w(x_1, \dots, x_d)$ be a non-trivial word, namely a non-identity element of the free group F_d on x_1, \dots, x_d .

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- Let G be a group. The word map $w : G^d \rightarrow G$ is defined by substituting group elements g_1, \dots, g_d in x_1, \dots, x_d respectively.
- Let $w(G) \subseteq G$ denote the image of this map, and denote

$$w(G)^k = \{g_1 \cdot g_2 \cdot \dots \cdot g_k \mid g_i \in w(G)\}.$$

Waring Type Problems in Finite Simple Groups

FSG = Finite (non-abelian) Simple Group.

Assume **CFSG** (the Classification).

Theorem (Wilson, 1994)

Any element of a FSG is a product of c commutators, where c is some absolute constant. I.e., for $w = [x, y] = x^{-1}y^{-1}xy$, $w(G)^c = G$. (c unspecified)

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Theorem (Martinez-Zelmanov, 1996, Saxl-Wilson, 1997)

Let $w = x^k$. There exist $f(k)$ such that either $w(G) = 1$ or $w(G)^{f(k)} = G$ for any FSG G .

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Question

*Are there extensions of these results to **general words** w ?*

Theorem (Liebeck-Shalev, 2001)

For *any word w* there exists a positive integer $c = c_w$ such that, for any FSG G , either $w(G) = 1$ or $w(G)^c = G$.

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Theorem (Shalev, 2009)

For *any* $w \neq 1$, there exists a positive integer $N = N_w$ such that

$$w(G)^3 = G$$

for every FSG G with $|G| \geq N$.

Proof uses **probabilistic methods** following Erdős

New proof by Nikolov-Pyber in 2011 using **Gowers' trick**.

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However, various words w are not surjective on all FSG, or even on almost all of them. E.g. x^n is not surjective whenever $(n, |G|) \neq 1$, so x^2 is **never surjective** on a FSG.

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Given $w \neq 1$, there exists a constant $N = N_w$ such that

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Corollary

Given $k \geq 1$ there exists N_k such that, if G is a FSG satisfying $|G| \geq N_k$, then every element of G is a product of two k -th powers.

Better solution to Waring problem in the non-commutative world!

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NO!

Example (Jambor-Liebeck-O'Brien, 2013)

$w = x^2 [x^{-2}, y^{-1}]^k$ is not surjective on $\mathrm{PSL}_2(q)$ for infinitely many q .

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Let w_n be the n -th Engel word. Then $w_n(G) = G$ for $G = \mathrm{PSL}_2(q)$ when $q \geq q_0(n)$.

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This, and computer experiments on other groups suggest:

Conjecture

Let G be a FSG. Let w_n be the n -th Engel word. Then $w_n(G) = G$.

Notation

For a conjugacy class C , we denote

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Theorem (Ellers-Gordeev, 1998)

*Thompson conjecture holds for **groups of Lie type** over a finite field F_q , provided $q > 8$.*

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Theorem (Shalev, 2008-2009)

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Denote by $U_{w(G)}$ the uniform distribution on $w(G)$.

Denote by $U_{w(G)}^{*k}$ the distribution induced on $w(G)^k$ by a k -fold product:

$$U_{w(G)}^{*k}(g) = \text{Prob} \left\{ g_1 g_2 \dots g_k = g \mid \begin{array}{l} g_1, \dots, g_k \text{ distribute uniformly} \\ \text{and independently in } w(G) \end{array} \right\}.$$

This is the distribution induced on G by a **k -step random walk on the (directed) Cayley graph of G with $w(G)$ as a generating set.**

Theorem (Larsen-Schul-Shalev, 2008-9)

For $w \neq 1$, $\left\| U_{w(G)}^{*2} - U_G \right\|_1 \rightarrow 0$ as $|G| \rightarrow \infty$ and G is FSG.

Thus the *mixing time* of the random walk on G with respect to $w(G)$ as a generating set is *2*.

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Corollary

If $w \neq 1$ then $\frac{|w(G)^2|}{|G|} \rightarrow 1$ as $|G| \rightarrow \infty$ (G FSG).

Another natural distribution induced on G by a word map:

$$P_{w,G}(g) = \text{Prob} \left\{ w(g_1, \dots, g_d) = g \mid \begin{array}{l} g_1 \dots g_d \text{ distribute uniformly} \\ \text{and independently in } G \end{array} \right\}.$$

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For any word $w \neq 1$ there exists $\epsilon = \epsilon_w > 0$ such that for all large FSG G and $g \in G$, we have $P_{w,G}(g) \leq |G|^{-\epsilon}$.

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$P_{w,G}$ is a **class function** on G , hence a linear combination of irreducible characters: $P_{w,G} = |G|^{-1} \sum_{\chi} a_{\chi} \chi$. (**Fourier expansion**)
Hence **character methods** are highly relevant.

Theorem (Garion-Shalev, 2009)

For FSG G , $\|P_{[x,y],G} - U_G\|_1 \rightarrow 0$ as $|G| \rightarrow \infty$.

Sketch of proof:

(i) $\|P_{[x,y],G} - U_G\|_1 \leq \sum_{\chi \neq 1} \chi(1)^{-2}$.

(ii) $\sum_{\chi \neq 1} \chi(1)^{-2} \rightarrow 0$ as $|G| \rightarrow \infty$.

Application to **Product Replacement Algorithm**. Similar result for x^2y^2 .

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Theorem (Larsen-Shalev, 2013)

Fix $n, m \geq 1$. Then for FSG G , $\|P_{x^m y^n, G} - U_G\|_1 \rightarrow 0$ as $|G| \rightarrow \infty$.

Work in progress: Same for $w_1 w_2$, where $w_1, w_2 \neq 1$ are words in disjoint variables.

For $w = x$ we have: $P_{x,G} \equiv \frac{1}{|G|}$ for every finite group G .

Theorem (Puder-Parzanchevski, 2011)

$P_{w,G} \equiv \frac{1}{|G|}$ for every finite group G if and only if w is a primitive word (there exists $\varphi \in \text{Aut}(F_d)$ with $\varphi(w) = x$).

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For any $w \in F_d$ and $\varphi \in \text{Aut}(F_d)$, $P_{w,G} \equiv P_{\varphi(w),G}$ for every finite group G .

(since $\varphi(w)(h_1, \dots, h_d) = w(g_1, \dots, g_d)$, where $h_i = \varphi^{-1}(x_i)(g_1, \dots, g_d)$).

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Question

Given $w, w' \in F_d$ such that $P_{w,G} = P_{w',G}$ for every finite group G , is there $\varphi \in \text{Aut}(F_d)$ with $\varphi(w) = w'$?

Extensions to Infinite Groups

G semisimple, simply connected, algebraic group over \mathbb{Q} .

Consider the p -adic group $G(\mathbb{Z}_p)$, and the arithmetic group $G(\mathbb{Z})$.

Can we extend results from finite simple groups to infinite p -adic and arithmetic groups?

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Condition on p necessary. Not true for $w(G(\mathbb{Z}_p))^2$.

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Condition on p necessary. Not true for $w(G(\mathbb{Z}_p))^2$.

Theorem (Avni-Gelander-Kassabov-Shalev, 2013)

If n is a proper divisor of $p - 1$ then every element of $\mathrm{PSL}_n(\mathbb{Z}_p)$ is a commutator.

Ore Conjecture for p -adic and arithmetic groups:

Question

*Suppose $n > 2$. Is every element of $\mathrm{SL}_n(\mathbb{Z}_p)$ a commutator?
Is every element of $\mathrm{SL}_n(\mathbb{Z})$ a commutator?*

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