On a Hamiltonian Problem For Triple Systems

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joint work with

V. Rödl, M. Schacht and E. Szemerédi

Erdős Centennial Conference

Theorem (Dirac 1952)

If an n-vertex graph G with $n \ge 3$ satisfies $\delta(G) \ge n/2$, then G is Hamiltonian.

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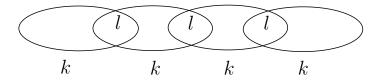
- What is a (Hamiltonian) cycle in a hypergraph?
- What replaces minimum degree in hypergraphs?

• *k*-uniform hypergraph H = (V, E), i.e., $E \subseteq \binom{V}{k}$

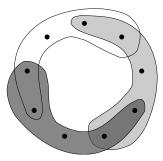
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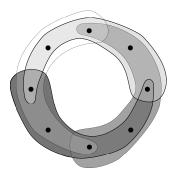
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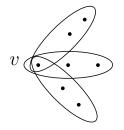
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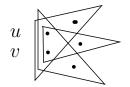


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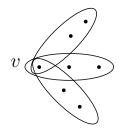
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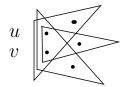
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• triple degrees $\delta_3(H)$

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Given integers k, ℓ , and d determine the function $h_d^{(k,\ell)}(n)$ with the property

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 Hamiltonian ℓ -cycle in H

for any *n*-vertex *k*-uniform hypergraph *H*.

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A. Rucínski (UAM Poznań & Emory)

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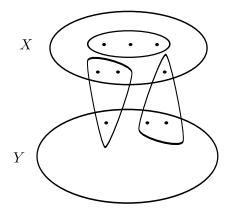
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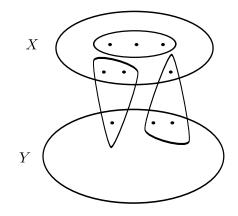
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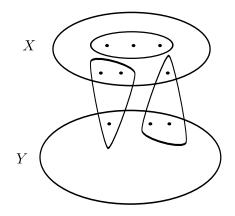


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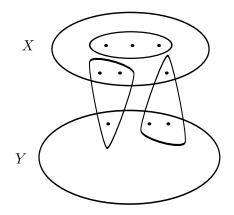
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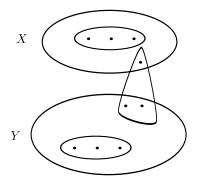
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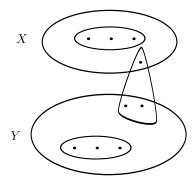
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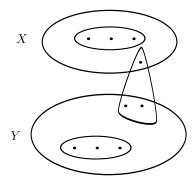
 $\Rightarrow \ \delta_1(H) \sim \binom{n}{2} - \binom{2n/3}{2} \sim \frac{5}{9}\binom{n}{2}$, but H contains no perfect matching



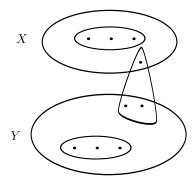
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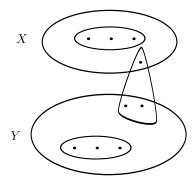
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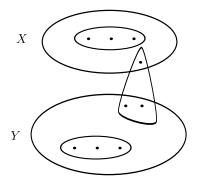
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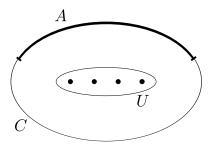
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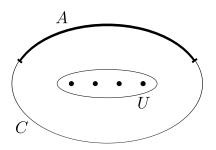
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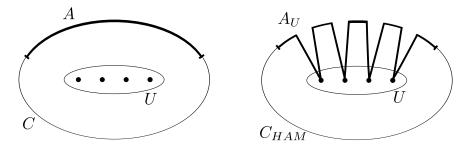
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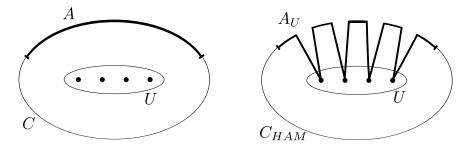
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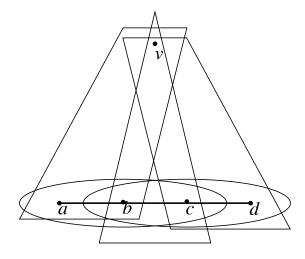
HAA...AM



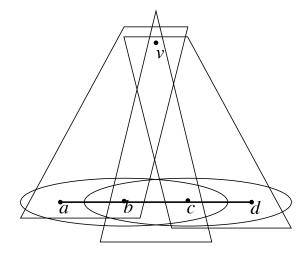
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Finding absorbers with $\delta_2(H) \ge (1/2 + \varepsilon)n$



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Fact: For every $v \in V$ there are $\varepsilon^2 n^4$ absorbers (a, b, c, d).

Absorbing path

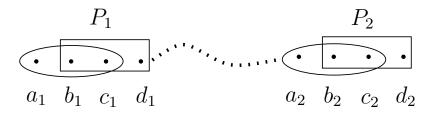
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- **3** Connect the selected 4-tuples P_i to obtain the path A



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- slightly more careful, remove only hyperedges which contain no pair of high degree
- balance between "finding absorbers" and "making connections between large pairs"

Questions



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