# On a Hamiltonian Problem For Triple Systems 

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## Erdős Centennial Conference

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If an n-vertex graph $G$ with $n \geq 3$ satisfies $\delta(G) \geq n / 2$, then $G$ is Hamiltonian.

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Problems:

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- What replaces minimum degree in hypergraphs?


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## Lower bound construction



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## HAA...AM



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Fact: For every $v \in V$ there are $\varepsilon^{2} n^{4}$ absorbers $(a, b, c, d)$.

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3 Connect the selected 4-tuples $P_{i}$ to obtain the path $A$


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■ balance between "finding absorbers" and "making connections between large pairs"

## Questions



