

On the digits of prime numbers

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Prime Number Theorem and Möbius Randomness Principle

p is always a prime number.

Von Mangoldt function: $\Lambda(n) = \log p$ if $n = p^k$, $\Lambda(n) = 0$ otherwise.

Prime Number Theorem (Hadamard, de la Vallée Poussin, 1896, indep.): $\sum_{n \leq x} \Lambda(n) = x + o(x)$.

Möbius function: $\mu(n) = (-1)^r$ if $n = p_1 \cdots p_r$ (distinct), $\mu(n) = 0$ if $\exists p, p^2 \mid n$.

Given a “reasonable” f , we say that f satisfies a PNT if we can get an asymptotic formula for $\sum_{n \leq x} \Lambda(n) f(n)$ while we say that f satisfies the MRP if $\sum_{n \leq x} \mu(n) f(n)$ is “small”.

These concepts are strongly related with Sarnak’s conjecture if f is produced by a zero topological entropy dynamical system.

For $f = 1$ these properties are equivalent: $\sum_{n \leq x} \mu(n) = o(x)$.

For more general f the MRP might be (slightly) less difficult to show than the PNT.

Gelfond's paper

In base $q \geq 2$ any $n \in \mathbb{N}$ can be written $n = \sum_{j \geq 0} \varepsilon_j(n) q^j$ where $\varepsilon_j(n) \in \{0, \dots, q-1\}$.

Theorem A (Gelfond, 1968) *The sum of digits $s(n) = \sum_{j \geq 0} \varepsilon_j(n)$ is well distributed in arithmetic progressions: given $m \geq 2$ with $(m, q-1) = 1$, there exists an explicit $\sigma_m > 0$ such that*

$$\forall m' \in \mathbb{N}^*, \forall (n', s) \in \mathbb{Z}^2, \sum_{\substack{n \leq x \\ n \equiv n' \pmod{m'} \\ s(n) \equiv s \pmod{m}}} 1 = \frac{x}{mm'} + O(x^{1-\sigma_m}).$$

Problem A (Gelfond, 1968)

1. Evaluate the number of prime numbers $p \leq x$ such that $s(p) \equiv a \pmod{m}$.
2. Evaluate the number of integers $n \leq x$ such that $s(P(n)) \equiv a \pmod{m}$, where P is a suitable polynomial [for example $P(n) = n^2$].

(Not so) Old results

Fouvry–Mauduit (1996):

$$\sum_{\substack{n \leq x \\ n=p \text{ or } n=p_1 p_2 \\ s(n) \equiv a \pmod{m}}} 1 \geq \frac{C(q, m)}{\log \log x} \sum_{\substack{n \leq x \\ n=p \text{ or } n=p_1 p_2}} 1.$$

Dartyge–Tenenbaum (2005): For $r \geq 2$,

$$\sum_{\substack{n \leq x \\ n=p_1 \dots p_r \\ s(n) \equiv a \pmod{m}}} 1 \geq \frac{C(q, m, r)}{\log \log x \log \log \log x} \sum_{\substack{n \leq x \\ n=p_1 \dots p_r}} 1.$$

Write $e(t) = \exp(2i\pi t)$.

Dartyge–Tenenbaum (2005) proved the Möbius Randomness Principle for $f(n) = e(\alpha s(n))$:

$$\sum_{n \leq x} \mu(n) e(\alpha s(n)) = O\left(\frac{x}{\log \log x}\right)$$

Sum of digits of primes

Theorem 1 (Mauduit-Rivat, 2010) *If $(q - 1)\alpha \in \mathbb{R} \setminus \mathbb{Z}$, there exists $C_q(\alpha) > 0$ and $\sigma_q(\alpha) > 0$,*

$$\left| \sum_{p \leq x} e(\alpha s(p)) \right| \leq C_q(\alpha) x^{1-\sigma_q(\alpha)}.$$

Corollary 1 *For $q \geq 2$ the sequence $(\alpha s(p_n))_{n \geq 1}$ is equidistributed modulo 1 if and only if $\alpha \in \mathbb{R} \setminus \mathbb{Q}$ (here $(p_n)_{n \geq 1}$ denotes the sequence of prime numbers).*

Corollary 2 *For $q \geq 2$, $m \geq 2$ such that $(m, q - 1) = 1$ and $a \in \mathbb{Z}$,*

$$\sum_{\substack{p \leq x \\ s(p) \equiv a \pmod{m}}} 1 \sim \frac{1}{m} \sum_{p \leq x} 1 \quad (x \rightarrow +\infty).$$

Theorem 2 (Drmota-Mauduit-Rivat, 2009) *local result: $s(p) = k$ for k “central”.*

Sum of digits of squares and polynomials

Theorem 3 (Mauduit-Rivat,2009) *If $(q - 1)\alpha \in \mathbb{R} \setminus \mathbb{Z}$, there exist $C_q(\alpha) > 0$ and $\sigma_q(\alpha) > 0$,*

$$\left| \sum_{n \leq x} e(\alpha s(n^2)) \right| \leq C_q(\alpha) x^{1-\sigma_q(\alpha)}.$$

Corollary 3 *For $q \geq 2$ the sequence $(\alpha s(n^2))_{n \geq 1}$ is equidistributed modulo 1 if and only if $\alpha \in \mathbb{R} \setminus \mathbb{Q}$.*

Corollary 4 *For $q \geq 2$, $m \geq 2$ such that $(m, q - 1) = 1$ and $a \in \mathbb{Z}$,*

$$\sum_{\substack{n \leq x \\ s(n^2) \equiv a \pmod{m}}} 1 \sim \frac{x}{m} \quad (x \rightarrow +\infty).$$

Theorem 4 (Drmota-Mauduit-Rivat,2011) *Idem for $s(P(n))$ where $P(X) \in \mathbb{Z}[X]$ is of degree $d \geq 2$, such that $P(\mathbb{N}) \subset \mathbb{N}$ and with leading coefficient a_d such that $(a_d, q) = 1$ and $q \geq q_0(d)$.*

Further questions

Are we able to extend these results to more general digital functions f ?

- For f strongly q -multiplicative, Martin-Mauduit-Rivat:

$$\left| \sum_{n \leq x} \Lambda(n) f(n) e(\theta n) \right| \leq C_q(f) x^{1-\sigma_q(f)}$$

- For block counting related functions (e.g. Rudin-Shapiro sequence)

$$f(n) = e \left(\alpha \sum_{j \geq 1} \varepsilon_{j-1}(n) \varepsilon_j(n) \right) ?$$

Wait and see...

Sum over prime numbers

By partial summation $\sum_{p \leq x} g(p) \longrightarrow \sum_{n \leq x} \Lambda(n) g(n)$ where $\Lambda(n)$ is von Mangoldt's function.

Advantage: convolutions !

$$\Lambda * \mathbb{1} = \log, \quad \text{i.e.} \quad \sum_{d|n} \Lambda(d) = \log n.$$

A classical process (Vinogradov, Vaughan, Heath-Brown) remains (using some more technical details), for some $0 < \beta_1 < 1/3$ and $1/2 < \beta_2 < 1$, to estimate uniformly the sums

$$S_I := \sum_{m \sim M} \left| \sum_{n \sim N} g(mn) \right| \quad \text{for } M \leq x^{\beta_1} \text{ (type I)}$$

where $MN = x$ (which implies that the “easy” sum over n is long) and for all complex numbers a_m, b_n with $|a_m| \leq 1, |b_n| \leq 1$ the sums

$$S_{II} := \sum_{m \sim M} \sum_{n \sim N} a_m b_n g(mn) \quad \text{for } x^{\beta_1} < M \leq x^{\beta_2} \text{ (type II),}$$

(which implies that both sums have a significant length).

Sums of type I

Key idea: the sum over n is free of unknown coefficients.

The knowledge of the function g should permit to estimate the sum $\sum_{n \sim N} g(mn)$.

In our case

$$g(mn) = f(mn) e(\theta mn)$$

where $f(n)$ is some digital function like $f(n) = e(\alpha s(n))$.

Some arguments from Fouvry and Mauduit (1996) can be generalized.

In particular θ easily disappears in the proof.

Sums of type II - Smoothing the sums

By Cauchy-Schwarz:

$$|S_{II}|^2 \leq M \sum_{m \sim M} \left| \sum_{n \sim N} b_n f(mn) e(\theta mn) \right|^2.$$

Expanding the square and exchanging the summations leads to a smooth sum over m , but also two free variables n_1 and n_2 with no control.

Van der Corput's inequality: for $z_1, \dots, z_L \in \mathbb{C}$ and $R \in \{1, \dots, L\}$,

$$\left| \sum_{\ell=1}^L z_\ell \right|^2 \leq \frac{L+R-1}{R} \left(\sum_{\ell=1}^L |z_\ell|^2 + 2 \sum_{r=1}^{R-1} \left(1 - \frac{r}{R}\right) \sum_{\ell=1}^{L-r} \Re(z_{\ell+r} \bar{z}_\ell) \right)$$

where $\Re(z)$ denotes the real part of z .

Now $n_1 = n + r$ and $n_2 = n$ so that the size of $n_1 - n_2 = r$ is small. It remains to estimate

$$\sum_{n \sim N} b_{n+r} \bar{b}_n \sum_{m \sim M} f(m(n+r)) \overline{f(mn)} e(\theta mr).$$

Carry propagation

If $f(n) = e(\alpha s(n))$, then in the difference $s(m(n+r)) - s(mn)$, the product mr is much smaller than mn . Take $M \asymp q^\mu$, $N \asymp q^\nu$ and $R \asymp q^\rho$ then

$$mn = \overbrace{35116790780999806546523475473462336857643565}^{\mu+\nu},$$

$$mr = \overbrace{396576345354568797095646467570}^{\mu+\rho},$$

we see that in the sum $mn+mr$ the digits after index $\mu+\rho$ may change only by carry propagation.

Proving that the number of pairs (m, n) for which the carry propagation exceeds

$$\mu_2 := \mu + 2\rho$$

is bounded by $O(q^{\mu+\nu-\rho})$, we can ignore them and replace $s(m(n+r)) - s(mn)$ by $s_{\mu_2}(m(n+r)) - s_{\mu_2}(mn)$ where s_{μ_2} is the truncated s function which considers only the digits of index $< \mu_2$:

$$s_{\mu_2}(n) := s(n \bmod q^{\mu_2})$$

which is periodic of period q^{μ_2} .

Sums of type II - Fourier analysis

We are now working modulo q^{μ_2} . For $f_{\mu_2}(n) = e(\alpha s_{\mu_2}(n))$ and its Discrete Fourier Transform

$$\widehat{f_{\mu_2}}(t) = \frac{1}{q^{\mu_2}} \sum_{0 \leq u < q^{\mu_2}} f_{\mu_2}(u) e\left(-\frac{ut}{q^{\mu_2}}\right).$$

By Fourier inversion formula and exchanges of summations we must show that the quantity

$$\sum_{0 \leq h < q^{\mu_2}} \sum_{0 \leq k < q^{\mu_2}} \left| \widehat{f_{\mu_2}}(h) \widehat{f_{\mu_2}}(-k) \right| \sum_{n \sim N} \left| \sum_{m \sim M} e\left(\frac{hm(n+r) + kmn}{q^{\mu_2}} + \theta mr\right) \right|$$

is estimated by $O(q^{\mu+\nu-\rho})$.

The geometric sum over m and the summation over n can be handled by classical arguments from analytic number theory. This can be done uniformly in θ .

The digital structure of f permits to prove the very strong L^1 estimate

$$\sum_{0 \leq h < q^{\mu_2}} \left| \widehat{f_{\mu_2}}(h) \right| = O(q^{\eta \mu_2}) \quad \text{with } \eta < 1/2.$$

This is sufficient to conclude for $f(n) = e(\alpha s(n))$.

The Rudin-Shapiro sequence

Let $f(n) = e\left(\frac{1}{2} \sum_{j \geq 1} \varepsilon_{j-1}(n) \varepsilon_j(n)\right) = (-1)^{\sum_{j \geq 1} \varepsilon_{j-1}(n) \varepsilon_j(n)}$.

$\widehat{f_{\mu_2}}$ is a Shapiro polynomial well known to have small absolute value: $\forall t \in \mathbb{R}, |\widehat{f_{\mu_2}}(t)| \leq 2^{\frac{1-\mu_2}{2}}$,
(with our normalization).

Pál Erdős always said that every talk should contain a proof. Let us study the L^1 norm of $\widehat{f_{\mu_2}}$.
From

$$1 = \sum_{0 \leq h < 2^{\mu_2}} |\widehat{f_{\mu_2}}(h)|^2 \leq 2^{\frac{1-\mu_2}{2}} \sum_{0 \leq h < 2^{\mu_2}} |\widehat{f_{\mu_2}}(h)|$$

we deduce

$$\sum_{0 \leq h < 2^{\mu_2}} |\widehat{f_{\mu_2}}(h)| \geq 2^{\frac{\mu_2-1}{2}}.$$

Therefore (so to say) $\eta = \frac{1}{2}$.

The proof for the sum of digits function cannot be adapted for the Rudin-Shapiro sequence.

A variant of van der Corput's inequality

(Introduced to solve Gelfond's problem for squares)

For $z_1, \dots, z_L \in \mathbb{C}$ and integers $k \geq 1$, $R \geq 1$ we have

$$\left| \sum_{\ell=1}^L z_\ell \right|^2 \leq \frac{L + kR - k}{R} \left(\sum_{\ell=1}^L |z_\ell|^2 + 2 \sum_{r=1}^{R-1} \left(1 - \frac{r}{R}\right) \sum_{\ell=1}^{L-kr} \Re(z_{\ell+kr} \overline{z_\ell}) \right).$$

For $k = 1$ this is the classical van der Corput's inequality.

Interest: control the indexes modulo k .

Taking $k = q^{\mu_1}$, this may permit to remove the lower digits.

Double truncation

Applying the classical Van der Corput inequality leads to replace f by

$$f_{\mu_2}(n) = e \left(\alpha \sum_{1 \leq j < \mu_2} \varepsilon_{j-1}(n) \varepsilon_j(n) \right).$$

Applying the variant of Van der Corput inequality with $k = q^{\mu_1}$ where $\mu_1 = \mu - 2\rho$ leads to replace f_{μ_2} by

$$f_{\mu_1, \mu_2}(n) = f_{\mu_2}(n) \overline{f_{\mu_1}(n)} = e \left(\alpha \sum_{\mu_1 \leq j < \mu_2} \varepsilon_{j-1}(n) \varepsilon_j(n) \right).$$

More generally we have proved that any digital function satisfying a carry propagation property can be replaced here by a function depending only on the digits of indexes $\mu_0, \dots, \mu_2 - 1$ for some μ_0 close to μ_1 , at the price of an acceptable error term.

For the Rudin-Shapiro sequence, $\mu_0 = \mu_1 - 1$.

Fourier analysis

After some technical steps the “digital part” and the “exponential sum part” are separated.

We need to estimate the following sum:

$$\sum_{|h_0| \leq q^{\mu_2 - \mu_0 + 2\rho}} \sum_{|h_1| \leq q^{\mu_2 - \mu_0 + 2\rho}} \min\left(\frac{q^{\mu_2 - \mu_0}}{\pi |h_0|}, 1\right) \min\left(\frac{q^{\mu_2 - \mu_0}}{\pi |h_1|}, 1\right) \\ \sum_{h_2 < q^{\mu_2 - \mu_0}} \sum_{h_3 < q^{\mu_2 - \mu_0}} |\hat{g}(h_0 - h_2) \hat{g}(h_3 - h_1) \hat{g}(-h_2) \hat{g}(h_3)| \\ \left| \sum_r \sum_s \left| \sum_m \sum_n e\left(\frac{(h_0 + h_1)mn + h_1mr + (h_2 + h_3)q^{\mu_1}sn}{q^{\mu_2}}\right) \right| \right|$$

where g is the $q^{\mu_2 - \mu_0}$ periodic function defined by

$$\forall k \in \mathbb{Z}, g(k) = f_{\mu_1, \mu_2}(q^{\mu_0}k).$$

End of the proof

The “exponential sum part” can be handled by appropriate estimates of exponential sums and similar tools. We need to average over all variables m, n, r, s .

For \hat{g} we need “only” that the L^∞ -norm is small.

This property is known for the sum of digits function and also for the classical Rudin-Shapiro sequence.

For generalized Rudin-Shapiro sequences we can prove it using a well chosen matrix norm.

In general it is very difficult.

Conclusion

We obtain a PNT and MRP for the Rudin-Shapiro sequence and its natural generalizations: counting $1 \underbrace{* \cdots *}_k 1$ for any $k \geq 0$, counting (overlapping) blocks $\underbrace{1 \cdots 1}_k$ for $k \geq 2$.

More generally we get a PNT and a MRP for any digital function satisfying a carry propagation property for which we can control uniformly the Discrete Fourier Transform.

General result – Definitions

Let $\mathbb{U} = \{z \in \mathbb{C}, |z| = 1\}$.

Definition 1 A function $f : \mathbb{N} \rightarrow \mathbb{U}$ has the carry property if, uniformly for $(\lambda, \kappa, \rho) \in \mathbb{N}^3$ with $\rho < \lambda$, the number of integers $0 \leq \ell < q^\lambda$ such that there exists $(k_1, k_2) \in \{0, \dots, q^\kappa - 1\}^2$ with

$$f(\ell q^\kappa + k_1 + k_2) \overline{f(\ell q^\kappa + k_1)} \neq f_{\kappa+\rho}(\ell q^\kappa + k_1 + k_2) \overline{f_{\kappa+\rho}(\ell q^\kappa + k_1)}$$

is at most $O(q^{\lambda-\rho})$ where the implied constant may depend only on q and f .

Definition 2 Given a non decreasing function $\gamma : \mathbb{R} \rightarrow \mathbb{R}$ satisfying $\lim_{\lambda \rightarrow +\infty} \gamma(\lambda) = +\infty$ and $c > 0$ we denote by $\mathcal{F}_{\gamma, c}$ the set of functions $f : \mathbb{N} \rightarrow \mathbb{U}$ such that for $(\kappa, \lambda) \in \mathbb{N}^2$ with $\kappa \leq c\lambda$ and $t \in \mathbb{R}$:

$$\left| q^{-\lambda} \sum_{0 \leq u < q^\lambda} f(uq^\kappa) e(-ut) \right| \leq q^{-\gamma(\lambda)}.$$

General result

Theorem 5 Let $\gamma : \mathbb{R} \rightarrow \mathbb{R}$ be a non decreasing function satisfying $\lim_{\lambda \rightarrow +\infty} \gamma(\lambda) = +\infty$, $c \geq 10$ and $f : \mathbb{N} \rightarrow \mathbb{U}$ be a function satisfying Definition 1 and $f \in \mathcal{F}_{\gamma, c}$ in Definition 2. Then for any $\theta \in \mathbb{R}$ we have

$$\left| \sum_{n \leq x} \Lambda(n) f(n) e(\theta n) \right| \ll c_1(q) (\log x)^{c_2(q)} x q^{-\gamma(2 \lfloor (\log x) / 80 \log q \rfloor) / 20},$$

with

$$c_1(q) = \max(\tau(q), \log^2 q)^{1/4} (\log q)^{-2 - \frac{1}{4} \max(\omega(q), 2)}$$

and

$$c_2(q) = \frac{9}{4} + \frac{1}{4} \max(\omega(q), 2).$$