

Infinite Sidon sequences

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The origin of the problem

In 1932, Simon Sidon asked to Erdős about the slowest possible growth of an infinite sequence A of positive integers having the property that all the sums

$$a + a', \quad a \leq a', \quad a, a' \in A$$

are distinct.

Erdős named them Sidon sequences and they became one of his favorite topics.

The Sidon sequence given by the greedy algorithm

Main problem: Construct (or prove the existence of) an infinite Sidon sequence A with counting function

$$A(x) = |A \cap [1, x]|$$

as large as possible.

Erdős considered the sequence given by the greedy algorithm:

- ▶ Starting with $a_1 = 1$, define a_{n+1} as the least positive integer we can add to the set $\{a_1, \dots, a_n\}$ preserving the Sidon property.

1, 2, 4, 8, 13, 21, 31, 45, 66, 81, 97, 123, 148, 182, 204, ...

The Sidon sequence given by the greedy algorithm

The forbidden elements for a_{n+1} are those of the form

$$a_i + a_j - a_k, \quad i, j, k \leq n.$$

Since there are at most n^3 of them, certainly we have that

$$a_{n+1} \leq n^3 + 1,$$

which implies

$$A(x) \gg x^{1/3}.$$

An upper bound for the counting function of a Sidon sequence

The number of sums $a + a'$, $a, a' \in A$ with $a < a' \leq x$ is $\binom{A(x)}{2}$, and all of them are distinct and less than $2x$:

$$\binom{A(x)}{2} < 2x \implies A(x) \ll x^{1/2}.$$

Conjecture (Erdős): For each $\epsilon > 0$ there is an infinite Sidon sequence A with

$$A(x) \gg x^{1/2-\epsilon}.$$

- ▶ Erdős proved that the conjecture is false for $\epsilon = 0$.

The construction of Ajtai, Komlós and Szemerédi

The greedy Sidon sequence found by Erdős was the densest known during almost 50 years.

Theorem (Ajtai, Komlós and Szemerédi, 1981): There exists an infinite Sidon sequence with counting function

$$A(x) \gg (x \log x)^{1/3}.$$

They wrote: *“The task of constructing a denser sequence has so far resisted all efforts, both constructive and random methods. Here we use a random construction for giving a sequence which is denser than the above trivial one”*.

The construction of Ruzsa

Theorem (Ruzsa, 1998): There exists an infinite Sidon sequence with counting function

$$A(x) = x^{\sqrt{2}-1+o(1)}.$$

The starting point of Ruzsa was the observation that the sequence of the prime numbers is a multiplicative Sidon sequence, or equivalently the sequence $(\log p)_{p \in P}$ is a Sidon sequence of real numbers.

The construction of Ruzsa

Ruzsa introduced a random parameter $\alpha \in [1, 2]$ and constructed a sequence $A_\alpha = (a_p)_{p \in P}$, where each integer a_p is built using the binary digits of $\alpha \log p$.

Then he proved that for almost all $\alpha \in [1, 2]$ it is possible to extract a dense Sidon sequence from A_α .

The constructions of Ruzsa and Ajtai, Komlós and Szemerédi are probabilistic. They are not explicit.

An explicit construction

It was an open problem to construct an explicit Sidon sequence with counting function $A(x) \gg x^c$ for some $c > 1/3$.

We construct an explicit Sidon sequence as dense as Ruzsa's sequence:

Theorem (C., 2012): There exists an infinite Sidon sequence A , **which can be explicitly constructed**, with counting function

$$A(x) = x^{\sqrt{2}-1+o(1)}.$$

Generalized basis

Given a sequence of positive integers

$$\bar{q} := 4q_1, \dots, 4q_j, \dots \quad (\text{the base}),$$

any positive integer a can be written, in only a way, in the form

$$a = x_1 + x_2(4q_1) + x_3(4q_1)(4q_2) + \dots + x_j(4q_1) \cdots (4q_{j-1}) + \dots$$

where the digits x_j satisfy

$$0 \leq x_j < 4q_j.$$

We represent the integer a in the form

$$a := \dots x_j \dots x_1.$$

Summing integers as vectors

If all the digits of a, a' satisfy

$$q_j < x_j, x'_j < 2q_j,$$

$$\begin{aligned} a &= x_{k_1} \dots \dots \dots x_1 \\ a' &= \quad \quad \quad x'_{k_2} \dots x'_1 \end{aligned}$$

we have

$$a + a' = (x_{k_1} + 0) \dots (x_{k_2+1} + 0) (x_{k_2} + x'_{k_2}) \dots (x_1 + x'_1).$$

Furthermore, the digits of $a + a'$ determine the lengths k_1, k_2 of a, a' , $a \leq a'$.

The construction: the base and the set of indexes

- 1) We consider a fix generalized base

$$\bar{q} := 4q_1, \dots, 4q_j, \dots,$$

where the q_j are primes satisfying

$$2^{2j-1} < q_j \leq 2^{2j}.$$

- 2) We use the set of the primes P as the indices

$$A = (a_p)_{p \in P}$$

and represent the elements a_p in the base \bar{q} as:

$$a_p = \dots x_j(p) \dots x_1(p).$$

The construction: the growth

- 3) Fix c , $0 < c < 1/2$ and make a partition of the set of the primes (the set of indices):

$$P = \bigcup_k P_k, \quad P_k = \{p : 2^{c(k-1)^2} < p \leq 2^{ck^2}\}.$$

Proposition: Assume that the elements a_p with $p \in P_k$ have exactly k digits in the base \bar{q} ,

$$a_p = x_k(p) \dots x_1(p).$$

Then we have

$$A_{\bar{q},c}(x) = x^{c+o(1)}.$$

The construction: the digits

4) For $p \in P_k$ we define the digits of

$$a_p := x_k(p) \dots x_1(p)$$

as follows: the digit $x_j(p)$ is given by the solution of

$$g_j^{x_j(p)} \equiv p \pmod{q_j}, \quad q_j < x_j(p) < 2q_j,$$

where g_j is a given generator of $\mathbb{F}_{q_j}^*$.

(The digit $x_j(p)$ is the discrete logarithm of p modulo q_j and it is unique modulo $q_j - 1$.)

- ▶ We will prove that if there is a repeated sum

$$a_{p_1} + a_{p_2} = a_{p'_1} + a_{p'_2}$$

then the primes involved, p_1, p_2, p'_1, p'_2 , must satisfy some relations.

- ▶ We will prove that these relations cannot hold if

$$c \leq \text{some value } c_0.$$

Lemma 1: If $a_{p_1} + a_{p_2} = a_{p'_1} + a_{p'_2}$ then there exist $k_2 \leq k_1$ such that

$$p_1, p'_1 \in P_{k_1}, \quad p_2, p'_2 \in P_{k_2}.$$

$$a_{p_1} = x_{k_1}(p_1) \cdots x_{k_2}(p_1) \cdots x_1(p_1)$$

$$a_{p_2} = x_{k_2}(p_2) \cdots x_1(p_2)$$

$$a_{p'_1} = x_{k_1}(p'_1) \cdots x_{k_2}(p'_1) \cdots x_1(p'_1)$$

$$a_{p'_2} = x_{k_2}(p'_2) \cdots x_1(p'_2)$$

Lemma 2: If $a_{p_1} + a_{p_2} = a_{p'_1} + a_{p'_2}$, $p_1, p'_1 \in P_{k_1}$, $p_2, p'_2 \in P_{k_2}$ then

$$p_1 p_2 \equiv p'_1 p'_2 \pmod{q_1 \cdots q_{k_2}}.$$

$$a_{p_1} = x_{k_1}(p_1) \cdots x_{k_2}(p_1) \cdots x_1(p_1)$$

$$a_{p_2} = \phantom{x_{k_1}(p_1) \cdots} x_{k_2}(p_2) \cdots x_1(p_2)$$

$$a_{p'_1} = x_{k_1}(p'_1) \cdots x_{k_2}(p'_1) \cdots x_1(p'_1)$$

$$a_{p'_2} = \phantom{x_{k_1}(p'_1) \cdots} x_{k_2}(p'_2) \cdots x_1(p'_2)$$

For $1 \leq j \leq k_2$ we have

$$x_j(p_1) + x_j(p_2) = x_j(p'_1) + x_j(p'_2)$$

$$g_j^{x_j(p_1) + x_j(p_2)} \equiv g_j^{x_j(p'_1) + x_j(p'_2)} \pmod{q_j}$$

$$p_1 p_2 \equiv p'_1 p'_2 \pmod{q_j}$$

Lemma 3: If $a_{p_1} + a_{p_2} = a_{p'_1} + a_{p'_2}$, $p_1, p'_1 \in P_{k_1}$, $p_2, p'_2 \in P_{k_2}$ then

$$p_1 \equiv p'_1 \pmod{q_{k_2+1} \cdots q_{k_1}}.$$

$$a_{p_1} = x_{k_1}(p_1) \cdots x_{k_2}(p_1) \cdots x_1(p_1)$$

$$a_{p_2} = x_{k_2}(p_2) \cdots x_1(p_2)$$

$$a_{p'_1} = x_{k_1}(p'_1) \cdots x_{k_2}(p'_1) \cdots x_1(p'_1)$$

$$a_{p'_2} = x_{k_2}(p'_2) \cdots x_1(p'_2)$$

For $k_2 + 1 \leq j \leq k_1$ we have

$$x_j(p_1) = x_j(p'_1)$$

$$g_j^{x_j(p_1)} \equiv g_j^{x_j(p'_1)} \pmod{q_j}$$

$$p_1 \equiv p'_1 \pmod{q_j}$$

$$\text{If } a_{p_1} + a_{p_2} = a_{p'_1} + a_{p'_2}$$

$$\text{i) } p_1, p'_1 \in P_{k_1} = \{p : 2^{c(k_1-1)^2} < p \leq 2^{ck_1^2}\} \quad (\text{Lemma 1})$$

$$p_2, p'_2 \in P_{k_2} = \{p : 2^{c(k_2-1)^2} < p \leq 2^{ck_2^2}\}$$

$$\text{ii) } p_1 p_2 \equiv p'_1 p'_2 \pmod{q_1 \cdots q_{k_2}} \quad (\text{Lemma 2})$$

$$\text{iii) } p_1 \equiv p'_1 \pmod{q_{k_2+1} \cdots q_{k_1}} \quad (\text{Lemma 3})$$

$$\text{iv) } 2^{2j-1} < q_j \leq 2^{2j} \quad (\text{by construction})$$

$$\begin{array}{ccc}
 \text{i)} & & \text{ii)} & & \text{iv)} \\
 \downarrow & & \downarrow & & \downarrow \\
 2^{ck_1^2 + ck_2^2} & \geq & |p_1 p_2 - p'_1 p'_2| & \geq & q_1 \cdots q_{k_2} > 2^{1+3+\cdots+(2k_2-1)} = 2^{k_2^2}
 \end{array}$$

$$\implies k_2^2 < \frac{c}{1-c} k_1^2.$$

$$\text{If } a_{p_1} + a_{p_2} = a_{p'_1} + a_{p'_2}$$

$$\text{i) } p_1, p'_1 \in P_{k_1} = \{p : 2^{c(k_1-1)^2} < p \leq 2^{ck_1^2}\} \quad (\text{Lemma 1})$$

$$p_2, p'_2 \in P_{k_2} = \{p : 2^{c(k_2-1)^2} < p \leq 2^{ck_2^2}\}$$

$$\text{ii) } p_1 p_2 \equiv p'_1 p'_2 \pmod{q_1 \cdots q_{k_2}} \quad (\text{Lemma 2})$$

$$\text{iii) } p_1 \equiv p'_1 \pmod{q_{k_2+1} \cdots q_{k_1}} \quad (\text{Lemma 3})$$

$$\text{iv) } 2^{2j-1} < q_j \leq 2^{2j} \quad (\text{by construction})$$

$$\begin{array}{ccc}
 \text{i)} & & \text{iii)} & & \text{iv)} \\
 \downarrow & & \downarrow & & \downarrow \\
 2^{ck_1^2} & \geq & |p_1 - p'_1| & \geq & q_{k_2+1} \cdots q_{k_1} > 2^{(2k_2+1)+\cdots+(2k_1-1)} = 2^{k_1^2 - k_2^2}
 \end{array}$$

$$\implies (1-c)k_1^2 < k_2^2.$$

An explicit infinite Sidon sequence

$$\begin{aligned}(1 - c)k_1^2 < k_2^2 < \frac{c}{1 - c}k_1^2 &\implies 1 - c < \frac{c}{1 - c} \\ &\implies c > \frac{3 - \sqrt{5}}{2} = 0.381966..\end{aligned}$$

Corollary (C., 2012): The sequence $A_{\bar{q},c}$ is a Sidon sequence for $c = \frac{3 - \sqrt{5}}{2} = 0.3819..$ with counting function

$$A_{\bar{q},c}(x) = x^{\frac{3 - \sqrt{5}}{2} + o(1)}.$$

An explicit infinite Sidon sequence

If $c > \frac{3-\sqrt{5}}{2}$ then $A_{\bar{q},c}$ is not a Sidon sequence. Some repeated sums $a_{p_1} + a_{p_2} = a_{p'_1} + a_{p'_2}$ may appear.

We remove the bad a_{p_1} involved in these sums to obtain a Sidon sequence.

If $c \leq \sqrt{2} - 1$, the removed elements are not too many.

Theorem (C., 2012): For $c = \sqrt{2} - 1$ the sequence $A_{\bar{q},c}$ contains a Sidon sequence A , **which can be explicitly constructed**, with

$$A(x) = x^{\sqrt{2}-1+o(1)}.$$

If $a_{p_1} + a_{p_2} = a_{p'_1} + a_{p'_2}$

$$p_1(p_2 - p'_2) = \frac{p_1 p_2 - p'_1 p'_2}{Q_1} \cdot Q_1 + \frac{(p'_1 - p_1) p'_2}{Q_2} \cdot Q_2,$$

where $Q_1 = q_1 \cdots q_{k_2}$, $Q_2 = q_{k_2+1} \cdots q_{k_1}$.

$$p_1(p_2 - p'_2) \in \left\{ s_1 \cdot Q_1 + s_2 \cdot Q_2 : |s_1| < \frac{2^{c(k_1^2 + k_2^2)}}{Q_1}, |s_2| \leq \frac{2^{c(k_1^2 + k_2^2)}}{Q_2} \right\}$$

for some k_2 , $k_2^2 < \frac{c}{1-c} k_1^2$.

We remove from each P_{k_1} all the primes $p_1 \in P_{k_1}$ dividing some integer of these sets.

It can be checked easily that the number of bad p_1 we have to remove is $o(|P_{k_1}|)$ for $c = \sqrt{2} - 1$.

B_h sequences

They are those sequences A such that all the sums of h elements of A are distinct. The greedy algorithm for B_h sequences gives one with

$$A(x) \gg x^{1/(2h-1)}.$$

Our approach also extends to B_h sequences:

Teorema (C., 2012) For each $h \geq 3$, **there exists** a B_h sequence A with

$$A(x) \gg x^{\sqrt{(h-1)^2+1}-(h-1)+o(1)}.$$

The cases $h = 3$ and $h = 4$ had been proved previously (C. and R, Tesoro, 2012) using a variant of Ruzsa's method, but that proof does not generalize to $h > 4$.

B_h sequences

For $h \geq 3$, our construction is not explicit. The problem is that we are not able to estimate the number of bad p_1 we have to remove from each P_{k_1} for a given basis \bar{q} .

We overcome this difficulty considering the probabilistic space of all basis

$$\bar{q} = h^2 q_1, \dots, h^2 q_j, \dots \quad \text{with} \quad 2^{2j-1} < q_j \leq 2^{2j}$$

and proving that for almost all basis \bar{q} the number of bad p_1 in each P_{k_1} is $o(|P_{k_1}|)$.

The finite Sidon set that motivated our construction

Let q be a prime and g a generator of \mathbb{F}_q^* and let $\log_g p$ be the discrete logarithm of p modulo q , which is unique modulo $q - 1$.

Theorem (C., 2012): The set

$$\mathcal{A} = \{\log_g p : p \text{ prime}, p \leq \sqrt{q}\}$$

is a Sidon set in \mathbb{Z}_{q-1} of size $\pi(\sqrt{q}) \sim \frac{\sqrt{q}}{\log \sqrt{q}}$.

$$\log_g p_1 + \log_g p_2 \equiv \log_g p'_1 + \log_g p'_2 \pmod{q-1}$$

$$p_1 p_2 \equiv p'_1 p'_2 \pmod{q}$$

$$p_1 p_2 = p'_1 p'_2$$