# Combinatorics, Probability and Computing 

http://journals.cambridge.org/CPC
Additional services for Combinatorics, Probability and
Computing:

Email alerts: Click here
Subscriptions: Click here
Commercial reprints: Click here
Terms of use : Click here

## A Note on Universal and Canonically Coloured Sequences

ANDRZEJ DUDEK, PETER FRANKL and VOJTĚCH RÖDL
Combinatorics, Probability and Computing / Volume 18 / Special Issue 05 / September 2009, pp 683-689 DOI: 10.1017/S0963548309009961, Published online: 21 May 2009

Link to this article: http://journals.cambridge.org/abstract S0963548309009961

## How to cite this article:

ANDRZEJ DUDEK, PETER FRANKL and VOJTĚCH RÖDL (2009). A Note on Universal and Canonically Coloured Sequences. Combinatorics, Probability and Computing, 18, pp 683-689 doi:10.1017/S0963548309009961

Request Permissions : Click here

# A Note on Universal and Canonically Coloured Sequences 

ANDRZEJ DUDEK ${ }^{1}$, PETER FRANKL ${ }^{2}$ and VOJTĚCH RÖDL ${ }^{3 \dagger}$<br>${ }^{1}$ Department of Mathematical Sciences, Carnegie Mellon University, Pittsburgh, PA 15213, USA (e-mail: adudek@andrew.cmu.edu)<br>${ }^{2}$ Peter Frankl Office, Ltd., 3-12-25 Shibuya, Shibuya-ku, Tokyo 150-0002, Japan<br>(e-mail: Peter111F@aol.com)<br>${ }^{3}$ Department of Mathematics and Computer Science, Emory University, Atlanta, GA 30322, USA<br>(e-mail: rodl@mathcs.emory.edu)

Received 19 June 2008; revised 12 March 2009; first published online 22 May 2009


#### Abstract

A sequence $X=\left\{x_{i}\right\}_{i=1}^{n}$ over an alphabet containing $t$ symbols is $t$-universal if every permutation of those symbols is contained as a subsequence. Kleitman and Kwiatkowski showed that the minimum length of a $t$-universal sequence is $(1-o(1)) t^{2}$. In this note we address a related Ramsey-type problem. We say that an $r$-colouring $\chi$ of the sequence $X$ is canonical if $\chi\left(x_{i}\right)=\chi\left(x_{j}\right)$ whenever $x_{i}=x_{j}$. We prove that for any fixed $t$ the length of the shortest sequence over an alphabet of size $t$, which has the property that every $r$-colouring of its entries contains a $t$-universal and canonically coloured subsequence, is at most $\mathrm{cr}^{\left\lfloor\frac{t}{2}\right\rfloor}$. This is best possible up to a multiplicative constant $c$ independent of $r$.


## 1. Introduction

A sequence $X=\left\{x_{i}\right\}_{i=1}^{n}$ over thealphabet $A=\left\{a_{1}, a_{2}, \ldots, a_{t}\right\}$ is $t$-universal if $X$ has as subsequences all permutations of the set $A$. For instance, if $A=\{1,2,3\}$, then 1231231 is 3 -universal. In general, the minimum length of $t$-universal sequences over an alphabet of size $t$, denoted by $f(t)$, is still unknown. The best-known upper bound is $f(t) \leqslant t^{2}-2 t+4$ for every $t \geqslant 3$, which was provided by several people (see, e.g., [2, 3, 4]). Moreover, Kleitman and Kwiatkowski [1] showed that $f(t)=(1-o(1)) t^{2}$.

In this note we consider the following Ramsey-type problem. We say that an $r$ colouring $\chi$ of the sequence $X=\left\{x_{i}\right\}_{i=1}^{n}$ is canonical if $\chi\left(x_{i}\right)=\chi\left(x_{j}\right)$ whenever $x_{i}=x_{j}$, i.e., all entries with the same value share the same colour. Let $\mathcal{R}(r, t)$ be the family of canonical Ramsey sequences $X$ over an alphabet of size $t$, i.e., sequences such that for

[^0]every $r$-colouring of the entries of $X$ there exists a $t$-universal and canonically coloured subsequence. Moreover, let
$$
f(r, t)=\min \{|X|: X \in \mathcal{R}(r, t)\} .
$$

Note that the number $f(r, t)$ is well-defined, i.e., $f(r, t)<\infty$. Indeed, let $X$ be a sequence over the alphabet $\left\{a_{1}, a_{2}, \ldots, a_{t}\right\}$ which consists of $(t-1) r^{t}+1$ consecutive blocks of the form $a_{1} a_{2} \cdots a_{t}$. Since there are exactly $r^{t}$ different ways to colour all entries of one particular block, at least $t$ blocks must have the same colour pattern. Clearly, the subsequence consisting of those $t$ blocks is $t$-universal and its colouring is canonical. We have just shown that $f(r, t) \leqslant\left((t-1) r^{t}+1\right) t=\mathcal{O}_{t}\left(r^{t}\right)$. The main result of this note determines the order of magnitude of $f(r, t)$ for a fixed integer $t$.

Theorem 1.1. For every positive integer $t$ there is a constant $c=c(t)$ such that for any $r$ the following inequalities hold:

$$
r^{\left\lfloor\frac{t}{2}\right\rfloor} \leqslant f(r, t) \leqslant c r^{\left\lfloor\frac{1}{2}\right\rfloor}
$$

Remark 1. We note that our proof of the lower bound yields a slightly stronger result. Namely, there exist two permutations $\sigma_{1}$ and $\sigma_{2}$ of the set $A$ of size $t$ such that any sequence over the alphabet $A$ and of length at most $r^{\left\lfloor\frac{t}{2}\right\rfloor}$ can be $r$-coloured in such a way that there is no canonically coloured subsequence containing $\sigma_{1}$ and $\sigma_{2}$.

## 2. Proof of Theorem 1.1

We will show that for a fixed $\ell$ there exists a constant $c=(2 \ell+1)(4 \ell+3)^{\ell}$ such that

$$
\begin{equation*}
\underbrace{r^{\ell}<f(r, 2 \ell)}_{(\mathrm{LB})} \leqslant \underbrace{f(r, 2 \ell+1) \leqslant c r^{\ell}}_{(\mathrm{UB})}, \tag{2.1}
\end{equation*}
$$

for any number of colours $r$. Clearly, this will imply Theorem 1.1. Note that since the second inequality holds trivially, we need to show (LB) and (UB) only.

### 2.1. The lower bound

In order to prove the lower bound (LB) we need to show that there is no sequence $X \in \mathcal{R}(r, 2 \ell)$ which has length $r^{\ell}$. To this end, we define an auxiliary sequence $U_{r, \ell}$ over an alphabet of size $2 \ell$, which contains all sequences of length $r^{\ell}$, and find an $r$-colouring of $U_{r, \ell}$ containing no $2 \ell$-universal and canonically coloured subsequence. Let $U_{r, \ell}$ be a sequence over the alphabet $A=\left\{a_{1}, a_{2}, \ldots, a_{2 \ell}\right\}$ consisting of $r^{\ell}$ consecutive blocks of the form $a_{1} a_{2} \cdots a_{2 \ell}$, i.e., $U_{r, \ell}=B^{(0)} B^{(1)} \cdots B^{\left(r^{\ell}-1\right)}$, where $B^{(i)}=x_{1}^{i} x_{2}^{i} \cdots x_{2 \ell}^{i}, x_{j}^{i}=a_{j}$ for any $0 \leqslant i \leqslant r^{\ell}-1$ and $1 \leqslant j \leqslant 2 \ell$. Observe that any sequence $X$ over the alphabet $A$ and of length $r^{\ell}$ is a subsequence of $U_{r, \ell}$. Hence, in order to show that $X \notin \mathcal{R}(r, 2 \ell)$ it is sufficient to show that $U_{r, \ell} \notin \mathcal{R}(r, 2 \ell)$. We are going to define an $r$-colouring $\chi_{r, \ell}$ of $U_{r, \ell}$ which has the property that there is no $2 \ell$-universal and canonically coloured subsequence in $U_{r, \ell}$.
Let $\chi_{r, \ell}: U_{r, \ell} \rightarrow\{0,1, \ldots, r-1\}$ be defined as follows. For a given integer $i, 0 \leqslant i \leqslant$ $r^{\ell}-1$, let $d_{\ell-1} d_{\ell-2} \cdots d_{0}$ be the $r$-nary expansion of $i$. Then, the $i$ th block of $U_{r, \ell}$ is
coloured as

$$
\begin{gather*}
\chi_{r, \ell}\left(x_{1}^{i}\right)=\chi_{r, \ell}\left(x_{2}^{i}\right)=d_{\ell-1} \\
\chi_{r, \ell}\left(x_{3}^{i}\right)=\chi_{r, \ell}\left(x_{4}^{i}\right)=d_{\ell-2} \\
\vdots  \tag{2.2}\\
\chi_{r, \ell}\left(x_{2 \ell-1}^{i}\right)=\chi_{r, \ell}\left(x_{2 \ell}^{i}\right)=d_{0} .
\end{gather*}
$$

For instance, if $\ell=1$, then $U_{r, 1}=a_{1} a_{2} a_{1} a_{2} \cdots a_{1} a_{2}$ is of length $2 r$. Set $q=r-1$. Then, $\chi_{r, 1}: U_{r, 1} \rightarrow\{0, \ldots, q\}$ gives on $U_{r, 1}$ the colour pattern $001122 \cdots q q$. Clearly, there is no canonically coloured subsequence which contains $a_{1} a_{2}$ and $a_{2} a_{1}$ as subsequences.

The next case $\ell=2$ illustrates the main idea of the general case. Let $\ell=2$. Then, $U_{r, 2}=B^{(0)} B^{(1)} \cdots B^{\left(r^{2}-1\right)}$, where $B^{(i)}=x_{1}^{i} x_{2}^{i} x_{3}^{i} x_{4}^{i}=a_{1} a_{2} a_{3} a_{4}$ for every $0 \leqslant i \leqslant r^{2}-1$. Set $q=r-1$. Below is the colour pattern induced by $\chi_{r, 2}$ :

$$
\begin{align*}
& 000000110022 \cdots 00 q q \\
& 110011111122 \cdots 11 q q \\
& 220022112222 \cdots 22 q q  \tag{2.3}\\
& q q 00 q q 11 \text { qq22 } \cdots \text { qqqq. }
\end{align*}
$$

Observe that in this colouring any subsequence of the form $a_{1} a_{2}$, more precisely, $x_{1}^{i} x_{2}^{j}$, $i \leqslant j$, has the property that

$$
\begin{equation*}
\chi_{r, 2}\left(x_{1}^{i}\right) \leqslant \chi_{r, 2}\left(x_{2}^{j}\right) . \tag{2.4}
\end{equation*}
$$

Also, for any subsequence $x_{2}^{i} x_{1}^{j}, i \leqslant j$, we have

$$
\begin{equation*}
\chi_{r, 2}\left(x_{2}^{i}\right) \leqslant \chi_{r, 2}\left(x_{1}^{j}\right) \tag{2.5}
\end{equation*}
$$

Now we show that there is no canonically coloured subsequence that contains $\sigma_{1}=$ $a_{1} a_{3} a_{4} a_{2}$ and $\sigma_{2}=a_{2} a_{4} a_{3} a_{1}$ as their subsequences. For a contradiction assume that this fails to be true. Since $x_{j}^{i}=a_{j}$ for all $0 \leqslant i \leqslant r^{2}-1$ and $1 \leqslant j \leqslant 4$, such $\sigma_{1}$ and $\sigma_{2}$ must be in $U_{r, 2}$ and be of the form

$$
x_{1}^{i_{1}} x_{3}^{i_{3}} x_{4}^{i_{4}} x_{2}^{i_{2}}=\sigma_{1}
$$

and

$$
x_{2}^{j_{2}} x_{4}^{j_{4}} x_{3}^{j_{3}} x_{1}^{j_{1}}=\sigma_{2}
$$

where

$$
\begin{equation*}
0 \leqslant i_{1} \leqslant i_{3} \leqslant i_{4} \leqslant i_{2} \leqslant r^{2}-1 \tag{2.6}
\end{equation*}
$$

and

$$
\begin{equation*}
0 \leqslant j_{2} \leqslant j_{4} \leqslant j_{3} \leqslant j_{1} \leqslant r^{2}-1 \tag{2.7}
\end{equation*}
$$

Moreover, due to our assumption, $\chi_{r, 2}\left(x_{1}^{i_{1}}\right)=\chi_{r, 2}\left(x_{1}^{j_{1}}\right), \chi_{r, 2}\left(x_{2}^{i_{2}}\right)=\chi_{r, 2}\left(x_{2}^{j_{2}}\right), \chi_{r, 2}\left(x_{3}^{i_{3}}\right)=\chi_{r, 2}\left(x_{3}^{j_{3}}\right)$ and $\chi_{r, 2}\left(x_{4}^{i_{4}}\right)=\chi_{r, 2}\left(x_{4}^{j_{4}}\right)$. This assumption together with (2.4) and (2.5) implies

$$
\chi_{r, 2}\left(x_{1}^{i_{1}}\right) \leqslant \chi_{r, 2}\left(x_{2}^{i_{2}}\right)=\chi_{r, 2}\left(x_{2}^{j_{2}}\right) \leqslant \chi_{r, 2}\left(x_{1}^{j_{1}}\right)=\chi_{r, 2}\left(x_{1}^{i_{1}}\right) .
$$

Consequently, $\chi_{r, 2}\left(x_{1}^{i_{1}}\right)=\chi_{r, 2}\left(x_{2}^{i_{2}}\right)=\chi_{r, 2}\left(x_{1}^{j_{1}}\right)=\chi_{r, 2}\left(x_{2}^{j_{2}}\right)$. That means that all indices $i_{1}, i_{2}$, $j_{1}$ and $j_{2}$ are in one row of (2.3), and so there exists an $m, 0 \leqslant m \leqslant r-1$, such that $m r \leqslant i_{1}, i_{2}, j_{1}, j_{2} \leqslant(m+1) r-1$. Consequently, by (2.6) and (2.7), $m r \leqslant i_{3}, i_{4}, j_{3}, j_{4} \leqslant(m+$ $1) r-1$ also holds. But then $\chi_{r, 2}\left(x_{3}^{i}\right) \leqslant \chi_{r, 2}\left(x_{4}^{i_{+}}\right)$and $\chi_{r, 2}\left(x_{4}^{j}\right)<\chi_{r, 2}\left(x_{3}^{j_{+}}\right)$for every $i \leqslant i_{+}$ and $j \leqslant j_{+}$such that $m r \leqslant i, i_{+}, j, j_{+} \leqslant(m+1) r-1$. In particular, $\chi_{r, 2}\left(x_{3}^{i_{3}}\right) \leqslant \chi_{r, 2}\left(x_{4}^{i_{4}}\right)=$ $\chi_{r, 2}\left(x_{4}^{j_{4}}\right)<\chi_{r, 2}\left(x_{3}^{j_{3}}\right)=\chi_{r, 2}\left(x_{3}^{i_{3}}\right)$, a contradiction.

Similarly, one can prove that for any $\ell>2$ there is no canonically coloured subsequence in $U_{r, \ell}$ with respect to $\chi_{r, \ell}(c f$. (2.2)) that contains both

$$
\begin{equation*}
\sigma_{\ell}^{1}=a_{1} a_{3} a_{5} \cdots a_{2 \ell-1} a_{2 \ell} \cdots a_{6} a_{4} a_{2} \tag{2.8}
\end{equation*}
$$

and

$$
\begin{equation*}
\sigma_{\ell}^{2}=a_{2} a_{4} a_{6} \cdots a_{2 \ell} a_{2 \ell-1} \cdots a_{5} a_{3} a_{1} \tag{2.9}
\end{equation*}
$$

as their subsequences. The proof goes by induction. Let us assume that $U_{r, \ell-1}$ has no canonically coloured subsequence in $\chi_{r, \ell-1}$ that contains both

$$
\sigma_{\ell-1}^{1}=a_{1} a_{3} a_{5} \cdots a_{2(\ell-1)-1} a_{2(\ell-1)} \cdots a_{6} a_{4} a_{2}
$$

and

$$
\sigma_{\ell-1}^{2}=a_{2} a_{4} a_{6} \cdots a_{2(\ell-1)} a_{2(\ell-1)-1} \cdots a_{5} a_{3} a_{1}
$$

i.e., $U_{r, \ell-1} \notin \mathcal{R}(r, 2(\ell-1))$. Suppose for a contradiction that $U_{r, \ell} \in \mathcal{R}(r, 2 \ell)$. In particular, there are indices

$$
\begin{equation*}
0 \leqslant i_{1} \leqslant i_{3} \leqslant i_{5} \leqslant \cdots \leqslant i_{2 \ell-1} \leqslant i_{2 \ell} \leqslant \cdots \leqslant i_{6} \leqslant i_{4} \leqslant i_{2} \leqslant 2 \ell \tag{2.10}
\end{equation*}
$$

and

$$
\begin{equation*}
0 \leqslant j_{2} \leqslant j_{4} \leqslant j_{6} \leqslant \cdots \leqslant j_{2 \ell} \leqslant j_{2 \ell-1} \leqslant \cdots \leqslant j_{5} \leqslant j_{3} \leqslant j_{1} \leqslant 2 \ell \tag{2.11}
\end{equation*}
$$

such that

$$
\begin{gathered}
x_{1}^{i_{1}} x_{3}^{i_{3}} x_{5}^{i_{5}} \cdots x_{2 \ell-1}^{i_{2 \ell-1}} x_{2 \ell}^{i_{2 \ell}} \cdots x_{6}^{i_{6}} x_{4}^{i_{4}} x_{2}^{i_{2}}=\sigma_{\ell}^{1}, \\
x_{2}^{j_{2}} x_{4}^{j_{4}} x_{6}^{j_{6}} \cdots x_{2 \ell}^{j_{2} \ell} x_{2 \ell-1}^{j_{2 \ell-1}} \cdots x_{5}^{j_{5}} x_{3}^{j_{3}} x_{1}^{j_{1}}=\sigma_{\ell}^{2},
\end{gathered}
$$

and $\chi_{r, \ell}\left(x_{k}^{i_{k}}\right)=\chi_{r, \ell}\left(x_{k}^{j_{k}}\right)$ for all $1 \leqslant k \leqslant 2 \ell$. As in the previous paragraph, one can prove that $\chi_{r, t}\left(x_{1}^{i_{1}}\right)=\chi_{r, t}\left(x_{2}^{i_{2}}\right)=\chi_{r, t}\left(x_{1}^{j_{1}}\right)=\chi_{r, t}\left(x_{2}^{j_{2}}\right)$. Then there exists an $m, 0 \leqslant m \leqslant r-1$, such that $m r^{\ell-1} \leqslant i_{1}, i_{2}, j_{1}, j_{2} \leqslant(m+1) r^{\ell-1}-1$, and consequently, by (2.10) and (2.11), also $m r^{\ell-1} \leqslant$ $i_{k}, j_{k} \leqslant(m+1) r^{\ell-1}-1$ for all $3 \leqslant k \leqslant 2 \ell$. Note that the subsequence $\tilde{U}$ of $U_{r, \ell}$ defined by elements $x_{k}^{i}$, $m r^{\ell-1} \leqslant i \leqslant(m+1) r^{\ell-1}-1,3 \leqslant k \leqslant 2 \ell$, is isomorphic to $U_{r, \ell-1}$. Moreover, the colouring $\chi_{r, \ell}$ restricted to $\tilde{U}$ corresponds to $\chi_{r, \ell-1}$. Hence, by induction, $\tilde{U}$ contains no canonically coloured subsequence containing both $\sigma_{\ell-1}^{1}$ and $\sigma_{\ell-1}^{2}$. Consequently, there is no canonically coloured subsequence in $U_{r, \ell}$ with $\sigma_{\ell}^{1}$ and $\sigma_{\ell}^{2}$, that is, $U_{r, \ell} \notin \mathcal{R}(r, 2 \ell)$.

### 2.2. The upper bound

In order to prove the upper bound (UB) we need to extend the concept of the universal sequences as follows. Let $t$ and $k, t \geqslant k$, be given integers. A variation of length $k$ on a set of size $t$ is a $k$-subset with a specific order. We say that a sequence over an alphabet of size $t$ is $(t, k)$-universal if every variation of length $k$ of those symbols is contained as a subsequence. For instance, the sequence 4123412314 is (4,3)-universal over the alphabet $\{1,2,3,4\}$. Let $\mathcal{R}(r, t, k)$ be the family of sequences $X$ over the alphabet of size $t$ with the property that, for every $r$-colouring of the entries of $X$, there exists a $(t, k)$-universal and canonically coloured subsequence. Moreover, let

$$
f(r, t, k)=\min \{|X|: X \in \mathcal{R}(r, t, k)\} .
$$

Note that $f(r, t)=f(r, t, t)$ and $f(r, t, 1)=t$.
First we show that

$$
\begin{equation*}
f(r, t, k+2) \leqslant(2 t r+1) f(r, t, k) \tag{2.12}
\end{equation*}
$$

for any $r \geqslant 1, t \geqslant 1, k \geqslant 1$ and $t \geqslant k+2$. Indeed, let $X \in \mathcal{R}(r, t, k)$ such that $|X|=f(r, t, k)$. Define a sequence $Y$ to be $2 t r+1$ consecutive copies of $X$, i.e., $Y=X^{(1)} X^{(2)} \cdots X^{(2 t r+1)}$, where $X^{(i)}=X$ for every $1 \leqslant i \leqslant 2 t r+1$. We show that $Y \in \mathcal{R}(r, t, k+2)$.

Fix a colouring $\chi: Y \rightarrow\{1,2, \ldots, r\}$. For a given symbol $a_{i}, 1 \leqslant i \leqslant t$, and colour $j$, $1 \leqslant j \leqslant r$, let $Y_{a_{i}, j}$ be the longest subsequence of $Y$ for which all entries are equal to $a_{i}$ and have the same colour $j$. Clearly $Y$ is a disjoint union over all $Y_{a_{i} j}$. For every $i \in\{1, \ldots, t\}$ and $j \in\{1, \ldots, r\}$, remove from $Y$ the first and last element of $Y_{a_{i}, j}$. Clearly, the total number of deleted entries is at most $2 t r$. Since $Y=X^{(1)} X^{(2)} \cdots X^{(2 t r+1)}$, there exists at least one copy of $X^{(i)}$ which is left untouched. But $X^{(i)} \in \mathcal{R}(r, t, k)$. Hence, there exists a $(t, k)$-universal and canonically coloured subsequence $\tilde{X}$ of $X^{(i)}$. Since we have already removed the endpoints of $Y_{a_{i} j}$, the sequence $\tilde{X}$ can be extended in $Y$ to a canonically coloured sequence $\tilde{Y}$ in which all symbols $\left\{a_{1}, \ldots, a_{t}\right\}$ appear before and also after $\tilde{X}$. This, together with $(t, k)$-universality of $\tilde{X}$, implies that every variation of length $k+2$ can be found in $\tilde{Y}$. In other words, $Y \in \mathcal{R}(r, t, k+2)$. Moreover, $|Y| \leqslant(2 t r+1) f(r, t, k)$, and hence (2.12) holds.

Applying (2.12) iteratively together with $f(r, t, 1)=t$ yields

$$
\begin{aligned}
f(r, t, 2 \ell+1) & \leqslant(2 t r+1) f(r, t, 2(\ell-1)+1) \\
& \leqslant(2 t r+1)^{\ell} f(r, t, 1)=(2 t r+1)^{\ell} t \leqslant t(2 t+1)^{\ell} r^{\ell} .
\end{aligned}
$$

Hence, in particular, $f(r, 2 \ell+1)=f(r, 2 \ell+1,2 \ell+1) \leqslant(2 \ell+1)(4 \ell+3)^{\ell} r^{\ell}$, which completes the proof of inequality (UB).

## 3. Concluding remarks

In the previous section we proved inequalities (2.1). It may be of interest to examine the behaviour of functions $f(r, 2 \ell)$ and $f(r, 2 \ell+1)$ in more detail. Below we propose the following problems.

Problem 3.1. Is it true that for a fixed $\ell$ we have

$$
\lim _{r \rightarrow \infty} \frac{f(r, 2 \ell)}{f(r, 2 \ell+1)}=1 ?
$$

Extending Problem 3.1, one can ask the following.
Problem 3.2. For a fixed $\ell$ determine the limits

$$
\lim _{r \rightarrow \infty} \frac{f(r, 2 \ell)}{r^{\ell}} \text { and } \lim _{r \rightarrow \infty} \frac{f(r, 2 \ell+1)}{r^{\ell}} .
$$

In Remark 1 we noted that the proof of the lower bound of $f(r, t)$ yields a stronger result. Following this remark, for a pair of permutations $\sigma_{1}$ and $\sigma_{2}$ of $[t]$, let us define $f_{\sigma_{1}, \sigma_{2}}(r, t)$ to be the length of shortest sequence over an alphabet of size $t$ which, for any $r$-colouring of its entries, contains a canonically coloured subsequence with both $\sigma_{1}$ and $\sigma_{2}$ as subsequences. For instance, we showed that for $\sigma_{\ell}^{1}$ and $\sigma_{\ell}^{2}(c f$. (2.8) and (2.9)), $r^{\ell} \leqslant f_{\sigma_{\ell}^{1}, \sigma_{\ell}^{2}}(r, 2 \ell) \leqslant f(r, 2 \ell)$.

Problem 3.3. Is it true that, for any permutations $\sigma_{1}$ and $\sigma_{2}$,

$$
\lim _{r \rightarrow \infty} \frac{f_{\sigma_{1}, \sigma_{2}}(r, t)}{f(r, t)}=0 ?
$$

In Problems 3.1-3.3 we assumed the size of the alphabet fixed and $r$ large. Swapping these assumptions, one might ask about the growth of $f(r, t)$ for fixed $r$ and $t$ large. For instance, for $r=1$, by [1], $f(1, t)=(1-o(1)) t^{2}$ holds. But even for $r=2$ we only know, by $(2.1)$, that $2^{\left\lfloor\frac{t}{2}\right\rfloor} \leqslant f(2, t) \leqslant t(2 t+1)^{\left\lfloor\frac{t}{2}\right\rfloor} 2^{\left\lfloor\frac{t}{2}\right\rfloor}$.

Finally, we consider the following related question. We say that two sequences $\left\{x_{i}\right\}_{i=1}^{n}$ and $\left\{y_{i}\right\}_{i=1}^{n}$ over integers are similar if their entries preserve the same order, i.e., $x_{i}<x_{j}$ if and only if $y_{i}<y_{j}$ for all $1 \leqslant i, j \leqslant n$. For a given sequence $X$ and an integer $r$, a sequence $Y$ is Ramsey if for every $r$-colouring of $Y$ there is a subsequence of $Y$ which is monochromatic and similar to $X$. Denote by $f(r, X)$ the length of the shortest Ramsey sequence $Y$. For instance, for two colours it is easy to see that $f(2, X) \leqslant|X|^{2}$. Indeed, let $X=\left\{x_{i}\right\}_{i=1}^{n}$ be a sequence over the alphabet $\{0, \ldots, n-1\}$. Then, note that the sequence $Y=Y^{(1)} Y^{(2)} \cdots Y^{(n)}$, where $Y^{(i)}=\left(n x_{i}+x_{1}, n x_{i}+x_{2}, \ldots, n x_{i}+x_{n}\right.$, for any $1 \leqslant i \leqslant n$, is Ramsey. Hence, $f(2, X) \leqslant|Y|=n^{2}$. On the other hand, one can also show that for $X=\left(1,2,3, \ldots,\left\lfloor\frac{n}{2}\right\rfloor, n, n-1, n-2, \ldots,\left\lfloor\frac{n}{2}\right\rfloor+1\right)$ every Ramsey sequence $Y$ has length at least $\frac{n^{2}}{4}$. Therefore, the bound $\mathcal{O}\left(|X|^{2}\right)$ on $f(2, X)$ is best possible. Extending the above construction one can prove that $f(r, X) \leqslant|X|^{r}$.

Problem 3.4. For a fixed $t$, estimate the order of magnitude of

$$
\max \{f(r, X): X \text { is a sequence over an alphabet of size } t\}
$$

as the function of $r$.

## References

[1] Kleitman, D. J. and Kwiatkowski, D. J. (1976) A lower bound on the length of a sequence containing all permutations as subsequences. J. Combin. Theory Ser. A 21 129-136.
[2] Koutas, P. J. and Hu, T. C. (1975) Shortest string containing all permutations. Discrete Math. 11 125-132.
[3] Mohanty, S. P. (1980) Shortest string containing all permutations. Discrete Math. 31 91-95.
[4] Newey, M. C. (1973) Notes on a problem involving permutations as subsequences. Computer Science Department Report, CS-73-340 Stanford University, Stanford, CA.


[^0]:    $\dagger$ Research partially supported by NSF grant DMS 0800070.

