

A new type of coding problem

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Let X be an n -element finite set, $0 < k < n/2$ an integer. Suppose that $\{A_1, B_1\}$ and $\{A_2, B_2\}$ are pairs of disjoint k -element subsets of X (that is, $|A_1| = |B_1| = |A_2| = |B_2| = k$, $A_1 \cap B_1 = \emptyset$, $A_2 \cap B_2 = \emptyset$). Define the distance of these pairs by $d(\{A_1, B_1\}, \{A_2, B_2\}) = \min\{|A_1 - A_2| + |B_1 - B_2|, |A_1 - B_2| + |B_1 - A_2|\}$. It is known ([2]) that this is really a distance on the space of such pairs and that the family of all k -element subsets of X can be paired (with one exception if their number is odd) in such a way that the distance of the pairs is at least k . Here we answer questions arising for distances larger than k .

1 Introduction

Let X be a finite set of n elements, $1 < k < n$ an integer. Unordered disjoint pairs $\{A, B\}$ of k -element sets (that is, $|A| = |B| = k$, $A \cap B = \emptyset$) will be considered. Define the *distance*

$$d(\{A_1, B_1\}, \{A_2, B_2\}) = \min\{|A_1 - A_2| + |B_1 - B_2|, |A_1 - B_2| + |B_1 - A_2|\}$$

between two such pairs. It has been verified in [2] that it is really a distance, that is, it satisfies the triangle inequality. We say that a set \mathcal{C} of such pairs is an (n, k, d) -code if the distance of any two elements is at least d .

Let $C(n, k, d)$ be the maximum size of an (n, k, d) -code. $C'(n, k, d)$ denotes the same under the additional condition that a

$$k\text{-element subset may occur only once in the pairs } \{A, B\} \in \mathcal{C} \text{ as } A \text{ or } B. \quad (1)$$

The following theorem was proved in [2].

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Theorem 1

$$C'(n, k, k) = \left\lfloor \frac{1}{2} \binom{n}{k} \right\rfloor.$$

It is obvious that one cannot choose more pairs using any k -element set at most once, so the theorem actually states that this many pairs can be constructed with pairwise distance k and satisfying (1). Theorem 1 is a sharpening of a theorem of [1] where $\lfloor \frac{1}{2} \binom{n}{k} \rfloor$ pairs were constructed under the condition that

$$\max\{|A_1 - A_2|, |B_1 - B_2|\}, \max\{|A_1 - B_2|, |B_1 - A_2|\} \geq \frac{k}{2}.$$

The method of the proofs of the constructions uses Hamiltonian type theorems.

It is quite natural to ask if one can choose $\lfloor \frac{1}{2} \binom{n}{k} \rfloor$ pairs with pairwise difference at least $k+1$. The answer is negative. In Section 2 we will give an upper estimate on $C(n, k, d)$ which will be less than $\lfloor \frac{1}{2} \binom{n}{k} \rfloor$ for $k < d$. Section 3 contains lower estimates on $C(n, k, d)$.

2 An upper estimate

Theorem 2 *Let $d \leq 2k \leq n$ be integers. Then*

$$C(n, k, d) \leq \frac{1}{2} \frac{n(n-1) \cdots (n-2k+d)}{k(k-1) \cdots \lceil \frac{d+1}{2} \rceil \cdot k(k-1) \cdots \lfloor \frac{d+1}{2} \rfloor}$$

holds.

PROOF: Let \mathcal{C} be a family of pairs of disjoint k -element subsets of X such that $d(C, C') \geq d$ for all $C, C' \in \mathcal{C}$ and count the number of pairs (C, D) where $C = \{A, B\} \in \mathcal{C}$, D is a $k - \lfloor \frac{d}{2} \rfloor$ -element subset of X and D is a subset of one of either A or B .

First, let us fix a $C = \{A, B\} \in \mathcal{C}$. There are exactly

$$2 \binom{k}{k - \lfloor \frac{d}{2} \rfloor} = 2 \binom{k}{\lfloor \frac{d}{2} \rfloor}$$

appropriate D s, therefore the total number of counted pairs (C, D) is

$$|\mathcal{C}| 2 \binom{k}{\lfloor \frac{d}{2} \rfloor}. \tag{2}$$

On the other hand, if D is fixed then suppose that $C_1 = \{A_1, B_1\}, C_2 = \{A_2, B_2\} \in \mathcal{C}$ and $D \subset A_1, A_2$. Since $|A_1 - A_2| \leq \lfloor \frac{d}{2} \rfloor$ therefore $|B_1 - B_2|$ must be at least $\lceil \frac{d}{2} \rceil$, that is, $|B_1 \cap B_2| \leq k - \lceil \frac{d}{2} \rceil$. Consequently the possible B s are subsets of the $n - k + \lfloor \frac{d}{2} \rfloor$ -element $X - D$ and they cannot cover the same $k - \lceil \frac{d}{2} \rceil + 1$ -element set. Hence the number of possible B s is at most

$$\frac{\binom{n-k+\lfloor \frac{d}{2} \rfloor}{k-\lceil \frac{d}{2} \rceil+1}}{\binom{k}{k-\lceil \frac{d}{2} \rceil+1}}.$$

The total number of pairs (C, D) cannot exceed

$$\binom{n}{k - \lfloor \frac{d}{2} \rfloor} \frac{\binom{n-k+\lfloor \frac{d}{2} \rfloor}{k-\lfloor \frac{d}{2} \rfloor+1}}{\binom{k}{k-\lfloor \frac{d}{2} \rfloor+1}}. \quad (3)$$

(2) \leq (3) leads to Theorem 2 by appropriate cancellations. \square

Corollary 3 *If $2 \leq k \leq n/2$ then*

$$C(n, k, k+1) < \lfloor \frac{1}{2} \binom{n}{k} \rfloor.$$

PROOF: Using Theorem 2 it is sufficient to prove

$$\frac{1}{2} \frac{n(n-1) \cdots (n-k+1)}{k(k-1) \cdots \lceil \frac{k+2}{2} \rceil \cdot k(k-1) \cdots \lfloor \frac{k+2}{2} \rfloor} < \frac{1}{2} \frac{n(n-1) \cdots (n-k+1)}{k!} - \frac{1}{2}. \quad (4)$$

It will be proved in the form

$$1 < \frac{n(n-1) \cdots (n-k+1)}{k!} \left(1 - \frac{\lfloor \frac{k}{2} \rfloor!}{k(k-1) \cdots (\lceil \frac{k}{2} \rceil + 1)} \right). \quad (5)$$

Observe that

$$2 \leq \frac{n}{k} < \frac{n-1}{k-1} < \cdots < \frac{n-k+1}{1}$$

and

$$\frac{1}{2} \geq \frac{\lfloor \frac{k}{2} \rfloor}{k} > \frac{\lfloor \frac{k}{2} \rfloor - 1}{k-1} > \cdots > \frac{1}{\lceil \frac{k}{2} \rceil + 1}.$$

Using these inequalities in (5) we arrive to the stronger inequality

$$1 < 2^k \left(1 - \frac{1}{2^{\lfloor \frac{k}{2} \rfloor}} \right)$$

which is trivially true for $2 \leq k$. (5), (4) and the corollary are proved. \square

3 Lower estimates

Let $1 < v < u < n$ be integers. The family \mathcal{P} is called an (n, u, v) *packing family* if it consists of u -element subsets of an n -element underlying set X and every v -element subset of X is contained in at most one member of \mathcal{P} . The class of all (n, u, v) packing families is denoted by $\mathbb{P}(n, u, v)$. Introduce the notation

$$m(n, u, v) = \max\{|\mathcal{P}| : \mathcal{P} \in \mathbb{P}(n, u, v)\}.$$

The inequality

$$m(n, u, v) \leq \frac{\binom{n}{v}}{\binom{u}{v}} \quad (1)$$

is obvious for a (n, u, v) packing family \mathcal{P} (1) holds with equality iff every v -element subset is contained in exactly one member of the family \mathcal{P} . In this case \mathcal{P} is called an (n, u, v) *Steiner family*.

The celebrated theorem of Rödl [5] (see also [3]) states that there are families asymptotically achieving the upper estimate (1), that is

$$\frac{m(n, u, v) \binom{u}{v}}{\binom{n}{v}} \rightarrow 1$$

for fixed u, v when n tends to infinity.

Proposition 4 *Let $d \leq 2k \leq n$ be integers. Then*

$$C(2k, k, d) \frac{n(n-1) \cdots (n-2k+d)}{(2k)(2k-1) \cdots d} (1+o(1)) \leq C(n, k, d) \quad (3)$$

holds, where $o(1)$ may depend on k and d .

PROOF: Take a family $\mathcal{P} \in \mathbb{P}(n, 2k, 2k-d+1)$ with size

$$|\mathcal{P}| = (1+o(1)) \frac{\binom{n}{2k-d+1}}{\binom{2k}{2k-d+1}} = (1+o(1)) \frac{n(n-1) \cdots (n-2k+d)}{(2k)(2k-1) \cdots d}$$

which exists by [5]. Let A_1, B_1 and A_2, B_2 be partitions (into k -element sets) of two different members of \mathcal{P} , that is, $A_1 \cap B_1 = A_2 \cap B_2 = \emptyset, |A_1| = |A_2| = |B_1| = |B_2| = k$ and $A_1 \cup B_1, A_2 \cup B_2$ are in \mathcal{P} . Their intersection has at most $2k-d$ elements, hence we have

$$|A_1 \cap A_2| + |B_1 \cap B_2|, |A_1 \cap B_2| + |B_1 \cap A_2| \leq |(A_1 \cup B_1) \cap (A_2 \cup B_2)| \leq 2k-d.$$

This implies

$$\begin{aligned} d(\{A_1, B_1\}, \{A_2, B_2\}) &= \min\{|A_1 - A_2| + |B_1 - B_2|, |A_1 - B_2| + |B_1 - A_2|\} = \\ &= \min\{k - |A_1 \cap A_2| + k - |B_1 \cap B_2|, k - |A_1 \cap B_2| + k - |B_1 \cap A_2|\} \geq d. \end{aligned}$$

Take the maximum number $C(2k, k, d)$ of such partitions with distance at least d in each member of \mathcal{P} . This construction proves (3). \square

Now we give a lower estimate on $C(2k, k, d)$ for some cases. The method is a modification of the method used by Sloane and Graham [4] proving lower bounds for constant weight codes.

Let us first consider the simplest case of $C(2k, k, 3) = C(2k, k, 4)$.

Theorem 5

$$|\mathcal{N}| \leq C(2k, k, 3)$$

where \mathcal{N} is the family of all k -element subsets A of $X = \{1, 2, \dots, 2k\}$ such that

$$\sum_{i \in A} i \equiv 0 \pmod{2k+1}.$$

PROOF: Since $1 + 2 + \dots + 2k = k(2k + 1) \equiv 0 \pmod{2k + 1}$ holds, $A \in \mathcal{N}$ implies $X - A \in \mathcal{N}$, too. \mathcal{N} consist of complementing pairs of k -element subsets of X .

Suppose that $A, B \in \mathcal{N}$, $|A \cap B| = k - 1$ holds.

$$\sum_{i \in A} i \equiv \sum_{i \in B} i \pmod{2k + 1}$$

implies

$$\sum_{i \in A - B} i \equiv \sum_{i \in B - A} i \pmod{2k + 1}.$$

Here $A - B$ and $B - A$ are 1-element sets, therefore they must be equal. Hence $A = B$, that is two different members of \mathcal{N} cannot have $k - 1$ common elements. They cannot have exactly one common element either, since this would imply that A and $X - B \in \mathcal{N}$ have $k - 1$ common elements, a contradiction. \square

It seems that $|\mathcal{N}|$ cannot be much smaller than

$$\frac{1}{(2k + 1)} \binom{2k}{k}.$$

We are quite sure that this is known, but we were unable to find the appropriate reference.

Suppose now that $q = 2k + 1$ is a prime power. We can prove an analogous lower bound for $C(2k, k, d)$ only in this case. Let $X = \{\omega_1, \dots, \omega_{q-1}\}$ be the set of all non-zero elements of the finite field $GF(q)$. Let $d = 2\delta$ and define $\mathcal{N}_0(k, \delta)$ as the family of all k -element subsets A of X such that

$$\sum_{i_1 < \dots < i_\rho \in A} \omega_{i_1} \cdots \omega_{i_\rho} = 0 \quad (4)$$

holds for every integer $1 \leq \rho < \delta$.

Let us see that $A \in \mathcal{N}_0(k, \delta)$ implies the same for $X - A$. Introduce the notation

$$s(B, u, v) = \sum \omega_j^u \omega_{i_1} \cdots \omega_{i_v}$$

for all $B \subset \{1, \dots, q - 1\}$, $0 \leq u, 0 \leq v < |B|$ where the sum is taken for all $v + 1$ different elements $j, i_1 < \dots < i_v$ of B . It is obvious that $s(B, 0, v)$ is $(|B| - v)$ times the sum of all products of v distinct ω_i s with indeces from B . On the other hand $s(B, 1, v) = \frac{(v+1)}{|B|-v} s(B, 0, v)$ holds and $s(B, u, 0)$ is the sum of the u th powers of ω_i s with indices from B .

$$s(B, u, 0) \frac{s(B, 0, v)}{|B| - v} = s(B, u, v) + s(B, u + 1, v - 1) \quad (1 \leq u, 1 \leq v < |B|) \quad (5)$$

is obviously true.

Let ε be a primitive root of the field. Then

$$s(X, u, 0) = \varepsilon^{0u} + \varepsilon^{1u} + \varepsilon^{2u} + \dots + \varepsilon^{(q-1)u} = \frac{\varepsilon^{qu} - 1}{\varepsilon^u - 1} = 0 \quad (6)$$

holds for $1 \leq u < q$.

(5) will be applied for $B = A$ several times. Start with the case $u = 1, v = \delta - 2$:

$$s(A, 1, 0) \frac{s(A, 0, \delta - 2)}{|A| - (\delta - 2)} = s(A, 1, \delta - 2) + s(A, 2, \delta - 3).$$

Here $s(A, 0, \delta - 2)$ and $s(A, 1, \delta - 2)$ are zero by (4). Consequently $s(A, 2, \delta - 3) = 0$ also holds. Applying (5) with $u = 2, v = \delta - 3$ and using $s(A, 2, \delta - 3) = 0$ the equality $s(A, 3, \delta - 4) = 0$ is obtained. Continuing this procedure we arrive to $s(A, \delta - 1, 0) = 0$. The equations $s(A, u, 0) = 0$ can be obtained in the same way for $1 \leq u \leq \delta - 1$. In other words,

$$\sum_{i \in A} \omega_i^u = 0 \quad (1 \leq u \leq \delta - 1) \quad (7)$$

holds. (6) and (7) imply that

$$s(X - A, u, 0) = s(X, u, 0) - s(A, u, 0) = \sum_{i \in X - A} \omega_i^u = 0$$

also holds for $1 \leq u \leq \delta - 1$. If the previous method is applied backwards for $X - A$, then it leads to the validity of (4) for $X - A$, proving that it is really in $\mathcal{N}_0(k, \delta)$.

We will now see that the symmetric difference of any two members A, B of $\mathcal{N}_0(k, \delta)$ is at least 2δ . Otherwise $A - B = \{r_1, \dots, r_\gamma\}, B - A = \{s_1, \dots, s_\gamma\}$ hold where $\gamma \leq \delta - 1$. Introduce the shorter notations $\alpha_i = \omega_{r_i}, \beta_i = \omega_{s_i}$. It is easy to see (see [4]) that the defining conditions (4) imply the equations

$$\begin{aligned} \sigma_1 &= \sum_i \alpha_i = \sum_i \beta_i, \\ \sigma_2 &= \sum_{i < j} \alpha_i \alpha_j = \sum_{i < j} \beta_i \beta_j, \\ &\dots \\ \sigma_{\delta-1} &= \sum_{i_1 < \dots < i_{\delta-1}} \alpha_{i_1} \dots \alpha_{i_{\delta-1}} = \sum_{i_1 < \dots < i_{\delta-1}} \beta_{i_1} \dots \beta_{i_{\delta-1}}. \end{aligned}$$

That is, the elementary symmetric functions of the α_i s and the β_i s agree, therefore $\alpha_1, \dots, \alpha_\gamma, \beta_1 \dots \beta_\gamma$ are all zeros of the polynomial

$$x^\gamma - \sigma_1 x^{\gamma-1} + \sigma_2 x^{\gamma-2} - \dots (-1)^\gamma \sigma_\gamma$$

of order γ . This contradiction proves that the pairwise distance of A and B is at least $d = 2\delta$. Since the same holds for the complements, the complementary pairs of the members of $\mathcal{N}_0(k, \delta)$ are really in distance at least d . The following theorem is proved.

Theorem 6 *If $2k + 1$ is a prime power and $d = 2\delta$ then*

$$\frac{1}{2} |\mathcal{N}_0(k, \delta)| \leq C(2k, k, d)$$

holds.

The size of $\mathcal{N}_0(k, \delta)$ can be determined for small values, but we believe that it cannot be much less than

$$\frac{1}{q^{\delta-1}} \binom{2k}{k},$$

since the defining sums are probably nearly equally distributed among all the $q^{\delta-1}$ possibilities.

4 Open problems

Theorem 2 and Proposition 4 imply the following statement.

Corollary 7

$$c_1(k, d)n^{2k-d+1} \leq C(n, k, d) \leq c_2(k, d)n^{2k-d+1}.$$

However we think that the upper bound of Theorem 2 is asymptotically correct.

Conjecture 8

$$\lim_{n \rightarrow \infty} \frac{C(n, k, d)}{n^{2k-d+1}} = \frac{1}{2k(k-1) \cdots \lceil \frac{d+1}{2} \rceil \cdot k(k-1) \cdots \lfloor \frac{d+1}{2} \rfloor}.$$

Actually we believe that, for an arbitrary pair of k and d , there are infinitely many n s with equality in Theorem 2.

The case $d = 1$ is uninteresting. If $d = 2$ then the upper and lower estimates coincide providing the $(n, 2k, 2k - 1)$ Steiner family exists. Therefore the first unfinished case is $d = 3$. Even in the case of $k = 2$, the upper and lower estimates significantly differ. The upper estimate is

$$C(n, 2, 3) \leq \frac{n(n-1)}{8}.$$

On the other hand $C(4, 2, 3)$ is obviously 1, therefore our construction gives only the lower bound

$$\frac{n(n-1)}{12}$$

when an $(n, 4, 2)$ Steiner family exists. Can one add $\frac{n(n-1)}{24}$ pairs of disjoint two-element sets (edges) to the Steiner system which preserves the condition that the pairwise distance of the pairs is at least 3?

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