

LARGE CLIQUES IN C_4 -FREE GRAPHS

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Dedicated to the memory of Paul Erdős

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A graph is called C_4 -free if it contains no cycle of length four as an induced subgraph. We prove that if a C_4 -free graph has n vertices and at least $c_1 n^2$ edges then it has a complete subgraph of $c_2 n$ vertices, where c_2 depends only on c_1 . We also give estimates on c_2 and show that a similar result does not hold for H -free graphs—unless H is an induced subgraph of C_4 . The best value of c_2 is determined for chordal graphs.

Graphs are understood to be simple, i.e. without loops or multiple edges and this is essential. The order of the largest complete subgraph of G is denoted by $\omega(G)$ and the order of the largest independent set of G is denoted by $\alpha(G)$. A graph is called C_4 -free if it contains no cycle of length four as an induced subgraph. The following question has been asked by Paul Erdős: is it true that C_4 -free graphs with n vertices and at least $c_1 n^2$ edges must contain complete subgraphs of $c_2 n$ vertices, where c_2 depends only on c_1 ? We give the affirmative answer ([Corollary 1](#)) with $c_2 = 0.4c_1^2$. In fact, we shall prove a more general result, [Theorem 1](#): A C_4 -free graph with n vertices and average degree at least a must contain a complete subgraph of order at least $0.1a^2 n^{-1}$. The role of C_4 is very important, similar results are not true for H -free graphs as the next proposition shows.

Proposition 1. *Suppose that H is a graph which is not an induced subgraph of C_4 . There are H -free graphs G_n with n vertices and at least $n^2/4$ edges with $\omega(G) = o(n)$.*

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Proof. If H contains three independent vertices then define G_n as a 'Ramsey graph', i.e. a graph on n vertices in which no three vertices form an independent set and $\omega(G_n)$ is as small as possible. It is well known that G_n has $n^2/2 - o(n^2)$ edges and $\omega(G_n) \leq c(n \log n)^{1/2} = o(n)$ ([4]). If $\alpha(H) \leq 2$ then one can define G_n as a balanced complete bipartite graph, in this case it has $\lfloor n^2/4 \rfloor$ edges and $\omega(G_n) = 2$. Since the induced subgraphs of G_n satisfying $\alpha(H) \leq 2$ are precisely the induced subgraphs of C_4 , the proof is finished.

Theorem 1. *Suppose that G is a C_4 -free graph on n vertices with average degree at least a . Then*

$$\omega(G) \geq 0.1a^2n^{-1}.$$

Applying [Theorem 1](#) for graphs with at least c_1n^2 edges we obtain the affirmative answer to the question of Erdős.

Corollary 1. *Let G be a C_4 -free graph on n vertices with at least c_1n^2 edges for some $0 < c_1 < \frac{1}{2}$. Then*

$$\omega(G) \geq 0.4c_1^2n.$$

The lower bound of $\omega(G)$ in [Corollary 1](#) is paralleled by two upper bounds. The first construction is better if c_1 is close to zero and the second is better if c_1 is close to $1/2$.

Construction 1. It is easy to check that clique substitutions into vertices of a graph which does not contain C_4 (not necessarily induced) yields C_4 -free graphs. It is well known that there are graphs on k vertices with approximately $(1/2)k^{3/2}$ edges which do not contain C_4 . These graphs are the so called 'polarity graphs' ([1], [3]). Substituting cliques into vertices of polarity graphs yields graphs on n vertices, at least c_1n^2 edges and with no cliques larger than $12c_1^2n$.

Construction 2. For two graphs, A and B , we denote by $A+B$ the graph C with vertex set $V(C) = V(A) \cup V(B)$ and edge set that consists of $E(A)$, $E(B)$ and all possible edges between $V(A)$ and $V(B)$.

Now consider a graph of form $A+B$ where A is a complete graph and B is a cycle with all chords of length less than $|B|/4$. When $1/4 \leq c_1 < 1/2$, with suitable sizes of A, B one can get C_4 -free graphs G with n vertices and at least c_1n^2 edges satisfying

$$\omega(G) \leq n \left(1 - \frac{3\sqrt{2}}{4} \sqrt{1 - 2c_1} \right).$$

Perhaps [Construction 2](#) is best possible for $c_1 \geq 1/4$. However we could not even decide the case $c_1 = 1/4$: is it true that a C_4 -free graph with n vertices and at least $n^2/4$ edges contains a clique of size $n/4$? (The best estimate we know is $n/6$, it comes by a special argument not shown in the paper.) A very special case we could not answer: assume that G is a $2k$ -regular C_4 -free graph with $4k+1$ vertices, is it true that $\omega(G) > k$? If true, it is best possible shown by the cycle C_{4k+1} with chords of length at most k .

Proof of Theorem 1. It is enough to prove [Theorem 1](#) for the case when the minimum degree of G is at least $a/2$. Indeed, it is well known that every graph G with average degree a contains an induced subgraph G^* with average degree $a^* \geq a$ such that the minimum degree in G^* is at least $a^*/2$ (see for example [\[2\]](#), Proposition 1.2.2). So if we know the theorem for G^* , we get

$$\omega(G) \geq \omega(G^*) \geq \frac{0.1(a^*)^2}{|V(G^*)|} \geq \frac{0.1a^2}{n}.$$

Assuming that the minimum degree of G is at least $a/2$, fix an independent set S with $|S| = t$. Set $S = \{x_1, x_2, \dots, x_t\}$. Let A_i be the set of neighbors of x_i in G and set $m = \max_{i \neq j} |A_i \cap A_j|$. Since G is C_4 -free, all the subgraphs $G[A_i \cap A_j]$ are complete graphs, and thus $m \leq \omega(G)$. Using that $|A_i| \geq a/2$ we get

$$ta/2 - \binom{t}{2}m \leq \sum_{1 \leq i \leq t} |A_i| - \sum_{1 \leq i < j \leq t} |A_i \cap A_j| \leq \left| \bigcup_{1 \leq i \leq t} A_i \right| \leq n$$

implying that

$$\omega(G) \geq m \geq \frac{ta/2 - n}{\binom{t}{2}}.$$

If $\alpha(G) \geq 4n/a$ then set $t = \lceil 4n/a \rceil$ and we get

$$\omega(G) \geq \frac{\lceil 4n/a \rceil a/2 - n}{\binom{\lceil 4n/a \rceil}{2}} \geq \frac{n}{\binom{\lceil 4n/a \rceil + 1}{2}}.$$

If $\alpha(G) \leq 4n/a$ then of course $\alpha(G) \leq \lfloor 4n/a \rfloor$ as well. Now we shall use the following claim: $\omega(G) \geq \frac{n}{\binom{\alpha(G)+1}{2}}$. This follows by selecting an independent set S with $|S| = \alpha(G) = \alpha$. Using the notation introduced above, the $\binom{\alpha}{2}$ sets $A_i \cap A_j$ and the α sets $\{x_i\} \cup B_i$ cover the vertex set of G where B_i denotes the set of vertices whose only neighbor in S is x_i . All of these sets span

complete subgraphs because G is C_4 -free and from the choice of S . Now we have

$$\omega(G) \geq \frac{n}{\binom{\alpha(G)+1}{2}} \geq \frac{n}{\binom{\lfloor 4n/a \rfloor + 1}{2}}.$$

Therefore, in both cases we have

$$\omega(G) \geq \frac{n}{\binom{\lfloor 4n/a \rfloor + 1}{2}} \geq \frac{n}{\binom{4n/a + 1}{2}} = \frac{a^2}{8n + 2a} \geq 0.1a^2/n$$

where in the last step, the trivial inequality $a \leq n$ was used for simple graphs. This finishes the proof.

Our next result shows that if c_1 tends to zero then [Corollary 1](#) can be improved by a factor of ten.

Theorem 2. *Assume that G is a C_4 -free graph with at least c_1n^2 edges. Then $\omega(G) \geq c_2^2n$ where c_2 tends to $4c_1^2$ if c_1 tends to zero.*

Proof. Let A denote the set of vertices in G whose degree is at least $10(c_1/c)^2n$ where c will be fixed later. Then

$$2c_1n^2 \leq 2e(G) = \sum_{x \in V} d(x) = \sum_{x \in A} d(x) + \sum_{x \in V-A} d(x) \leq \sum_{x \in A} d(x) + 10n^2(c_1/c)^2$$

which gives

$$\sum_{x \in A} d(x) \geq (2c_1 - 10(c_1/c)^2)n^2.$$

Case 1. There exists $x \in A$ whose neighborhood N in G span at least $c|N|^2$ edges. Now [Corollary 1](#) can be applied to $G[N]$ which gives a complete subgraph of at least $0.4c^2|N| \geq 0.4c^2 10(c_1/c)^2n = 4c_1^2n$ vertices proving the theorem.

Case 2. For every $x \in A$, the neighborhood of x , N_x , span at most $c|N_x|^2$ edges. This means that at least $\binom{d(x)}{2} - cd(x)^2$ edges are missing from $G[N_x]$. Let us denote by e the number of edges in G . Then one can estimate p , the number of induced paths with three vertices in G as

$$\begin{aligned} p &\geq \sum_{x \in A} \left(\binom{d(x)}{2} - cd(x)^2 \right) = \sum_{x \in A} \left(\frac{d(x)(d(x) - 1)}{2} - cd(x)^2 \right) = \\ &= (1/2 - c) \sum_{x \in A} d(x)^2 - e \geq (1/2 - c)n^{-1} \left(\sum_{x \in A} d(x) \right)^2 - e \geq \\ &\geq (1/2 - c)n^3 \left(2c_1 - 10(c_1/c)^2 \right) - e. \end{aligned}$$

Since the p paths connect at most $\binom{n}{2}$ pairs there are at least $p\binom{n}{2}^{-1}$ paths connecting the same vertex pair and the midpoints of these paths must form a complete graph since G is C_4 -free. Therefore

$$\omega(G) \geq (1 - 2c)(2c_1 - 10(c_1/c)^2)n - 1$$

which proves the theorem if c is selected so that both c and c_1/c tends to zero with c_1 .

A graph is called *chordal* if every cycle of length at least four contains a chord. Chordal graphs are clearly induced C_4 -free, thus [Theorem 1](#) and [Corollary 1](#) are valid for them. For chordal graphs we get (asymptotically) the best result.

Theorem 3. *Suppose that G is a chordal graph on n vertices with at least $\lceil c_1 n^2 \rceil$ edges, then*

$$\lfloor (1 - \sqrt{1 - 2c_1})n \rfloor \leq \omega(G).$$

Moreover, the the chordal graph which is the sum of an independent set of order $t = \lceil n\sqrt{1 - 2c_1} \rceil$ and a complete graph K_{n-t} shows that the inequality asymptotically is best possible.

Proof. It is well known that chordal graphs are perfect. See [2] Proposition 5.5.2. Thus the vertices of a chordal graph G can be partitioned into $\omega = \omega(G)$ independent sets $X_1, X_2, \dots, X_\omega$. Since G is chordal, the edges between X_i and X_j form an acyclic subgraph of G . Thus, we obtain

$$c_1 n^2 \leq \lceil c_1 n^2 \rceil \leq \sum_{i < j} (|V(X_i)| + |V(X_j)| - 1) = (\omega - 1)n - \binom{\omega}{2} \leq \omega n - \omega^2/2.$$

Thus $\omega^2 - 2\omega n + 2c_1 n^2 \leq 0$ from which the stated lower bound follows.

To check that the inequality is asymptotically sharp for the chordal graph defined in the theorem is routine so it is omitted.

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