

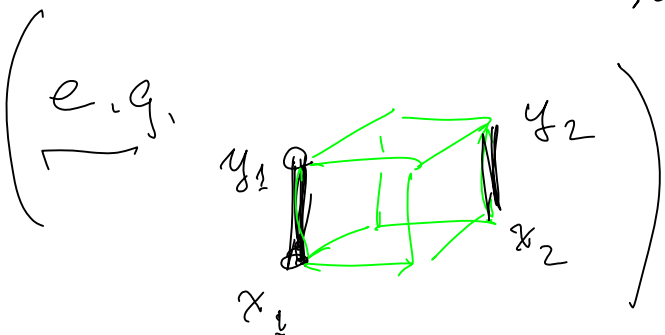
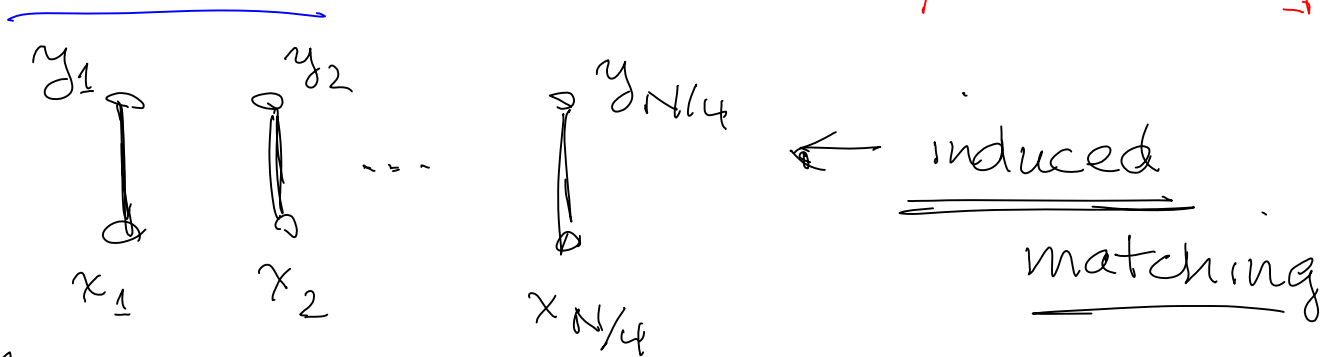
Q (Duffus-Frankl-Rödl):

How many max'l ind. sets in $\{0,1\}^d$?

$\dots Q^d$

lower bd: $(2d) 2^{N/4}$

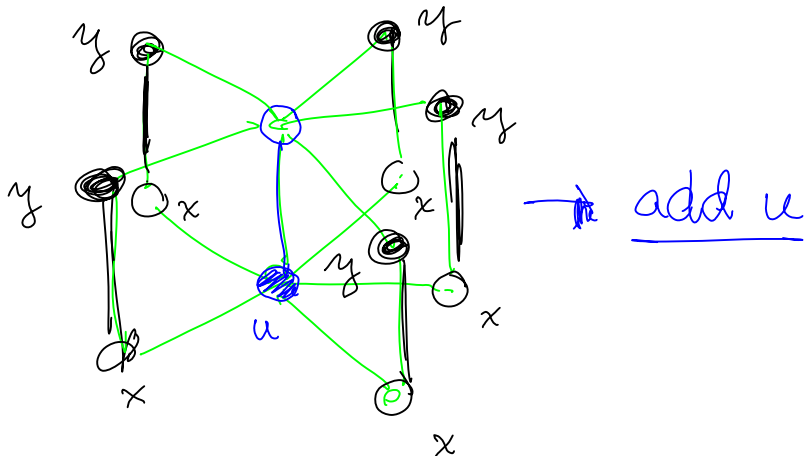
$[N := 2^d]$:



and do:

$$|I \cap \{x_i, y_i\}| = 1 \quad \forall i$$

add still-legals



$$2^{N/4} \text{ I's}$$

DFR:

Thm $N/4 < \log \text{m.i.s.}(Q^d) < 0.39 N$

$\rightarrow \# \dots$

Q: asymptotics of \log ?

(we'll come back)

Themes:

- ⊙ "sampling" \rightarrow (trivial?)
 - ⊙ entropy.
- } "INFO"
≅

E.g. of "sampling"

\mathcal{L}_k proj. pl. $|line| = k$ (order $k-1$)

τ (= # pts to meet all lines) = k

(any line is a k -cover)

Q (Erdős & Lovász 73)

How many random lines force $\tau = k$?

ANS (K92; as conj'd) $\Theta(k \log k)$

[l.b.: HW; why?: never mind]

(V. ROUGH:) $P = \{\text{points}\}$

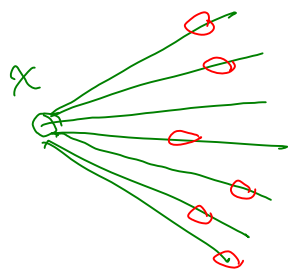
CRUX: bound

$|\{x \in \binom{P}{k-1} : \text{few lines miss } X\}|$

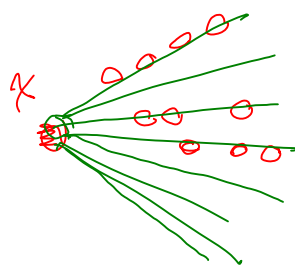
\mathbb{R}^2/x $\frac{1}{x}$, bound depends on x

X does $\textcircled{A} \Rightarrow$ most lines tangent to X

X from x :



$x \notin X$



$x \in X$



▶ "SAMPLE": \exists small $X_0 \subseteq X \Rightarrow$

$L(x, X_0)$ $\left\{ \begin{array}{l} \text{small for most } x \in X \\ \text{large for most } x \notin X \end{array} \right.$

$\left[\text{lines through } x \text{ meeting } X_0 \right]$

\rightarrow choose (cheap):

⊙ X_0

⊙ $x \in X \quad \bar{\omega} \quad L(x, X_0) \text{ large}$

⊙ $x \notin X \quad \bar{\omega} \quad L(x, X_0) \text{ small}$

$(\dots \text{ } \overline{\omega})$

"SKETCH" of X : $\{x : L(x, X_0) \text{ small}\}$

More e.q.'s

1 $i(G) := \#$ of indept sets in G

Thm G d -regular, N vertices

$$\Rightarrow i(G) \leq (2^{d+1} - 1)^{N/2d}$$

⊙ sharp for disjt $K_{d,d}$'s

⊙ suggested by Alon 91

(\leftrightarrow counting sum-free sets...)

⊙ G bipartite: $K_{0,1}$; entropy

⊙ gen'l: [REDACTED]

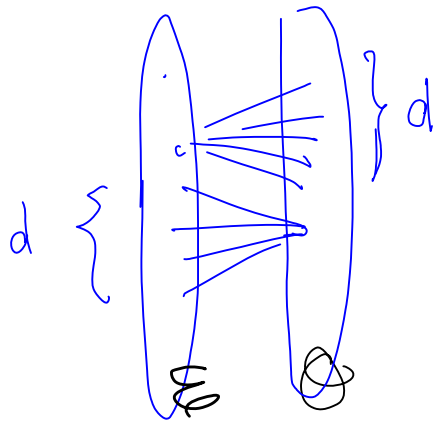
(hints: (1) reduce to bipartite

(2) a little embarrassing)

$$H(X) = \sum p(x) \log \frac{1}{p(x)}$$

$$(p(x) = \Pr(X=x))$$

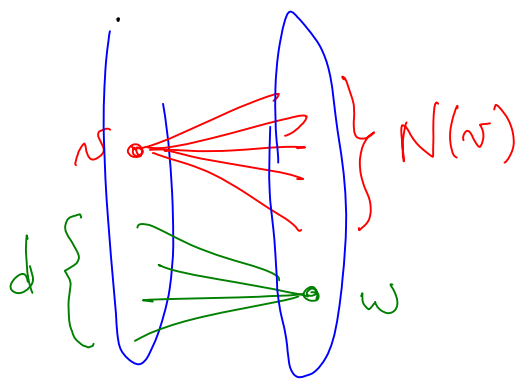
Entropy \leftrightarrow information $\left[\log_2(\#) = \underline{\underline{\text{info}}} \right]$



\mathbb{I} unif. ind.:

$$H(\mathbb{I}) = H(\mathbb{I} \cap \mathcal{O}) +$$

$$H(\mathbb{I} \cap \mathcal{E} | \mathbb{I} \cap \mathcal{O})$$



IDEA: info in one of

$$\sum_{\{v \in \mathbb{I}\}} \mathbb{I} \cap N(v)$$

MAIN: Shearer's Lemma \implies

$$H(\mathbb{I} \cap \mathcal{O}) \leq d^{-1} \sum_{v \in \mathcal{E}} H(\mathbb{I} \cap N(v))$$

$w \in \mathcal{O}$: " \sum " gives info on w d times

2

thm (Korshunov-Sapozhenko 83

⇒ Sapozhenko 89 ("sampling"))*

$$i(Q^d) \sim 2 \sqrt{e} 2^{2^{d-1}}$$

even/odd

"flaws"

all-even
ind sets



* exposition: D. Galvin [http ...](http://...)

$\boxed{3}$ $G = \mathbb{Z}_{2N}^d$ d large, fixed $\left[\begin{array}{l} \mathbb{R} \\ \mathbb{Z}^d \dots \end{array} \right]$
 bipartite \rightarrow N arb.

x, y even vertices $\left[\begin{array}{l} \text{think =} \\ \text{far apart.} \end{array} \right]$
 z odd vertex

Thm (Galvin - K 04)

\mathbb{I} unif. ind set:

$$\Pr(x \in \mathbb{I} | y \in \mathbb{I}) \approx 1/2$$

$$\Pr(x \in \mathbb{I} | z \in \mathbb{I}) \approx 0$$

\leftarrow pub date

\leftarrow we never got around to writing it date.

Thm (GKR 03, Peled 10)

σ unif 3-coloring (R, B, G):

$$\Pr(\sigma(x) = R | \sigma(y) = R) \approx 2/3$$

Conj: Sim. for k colors ($\exists d \geq d(k)$)

4 \mathcal{Y} unif $\in \{\text{homo's} : \mathbb{Q}^d \rightarrow \mathbb{Z} \text{ w } \underline{0} \rightarrow \uparrow 0\}$

Thm (K01)

entropy ($\# \dots$)

$$\exists b : \Pr(\underline{R(\mathcal{Y})} > b) < e^{-\Omega(d)}$$

$\approx |\text{range}(\mathcal{Y})|$

Conj (Benjamini-Häggström-Mossel)

$$\Pr(R(\mathcal{Y}) > \varepsilon n) \rightarrow 0 \quad \forall \varepsilon \downarrow$$

Thm (Galvin 03, conj'd K)

entropy \neq
"Sapozhenko"

$$\Pr(\underline{R(\mathcal{Y})} > \underline{s}) < e^{-\Omega(d)}$$

$$(\Pr(R(\mathcal{Y}) = s) > \Omega(1))$$

Conj (BHM) ~~\mathbb{Q}^d~~ n -vertex, bip. G

$\rightarrow \mathbb{E} R(\mathcal{Y})$ max for path
(e.g.)

Back to maximal indept sets

⌈ Nay, this is too much, to remember at night all the foolish things that were said in the morning.

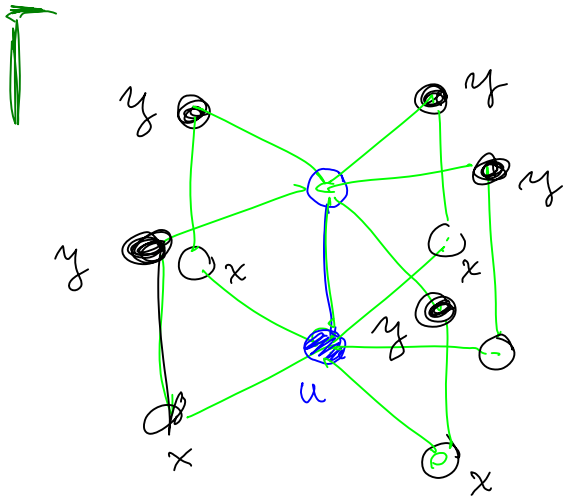
— Mr. Bingley ↓

Thm (Ilinca-K) G d -regular on N \Rightarrow

$$\log \text{m.i.s.}(G) < (1 + o_d(1)) \begin{cases} N/4 & \text{bipartite} \\ \frac{\log 3}{6} N & \text{gen'l} \end{cases}$$

[HW: conjecture worst examples]

[d-req, N vertices, bipartite]



← canon: IDEA
 $I \cap N(u)$ typ. large

step 1: $J \sim \text{Bin}(I, d^{-\epsilon}) \rightarrow$

(a) $|J| \approx d^{-\epsilon} N$ (small)

(b) $|\{x : |I \cap N(x)| > d^{2\epsilon}, x \notin J\}|$ small

→ cheap: specify J & these

step 2 $X = \{x : \text{still don't know}\}$

$\subseteq \{x : |I \cap N(x)| \leq d^{2\epsilon}\} \left[x \supseteq I \setminus J \right]$

step 2 $X = \{x : \text{still don't know}\}$:

$$\underline{|I \cap N(x)| \leq d^{2\varepsilon} \quad \forall x \in X} \quad \text{⊗}$$

$$\Upsilon = \{x \in X : d_x(y) > d^{3\varepsilon}\}$$

Π unif from $\{\underline{I}'_s \quad \bar{\omega} \quad \text{⊗}\} \rightarrow$

$$\log(\# \text{ of } \underline{I \cap \Upsilon}'_s) \leq H(\Pi \cap \Upsilon)$$

Shearer \leq $d^{-3\varepsilon} \sum_{x \in X} H(\Pi \cap N_\Upsilon(x))$

$\approx d^{2\varepsilon} \log d$ by ⊗

$$\leq (d^{-\varepsilon} \log d) N \quad \underline{\text{cheap}}$$

step 3 $Z = X \setminus Y$ (need $I \cap Z$)

\Rightarrow all degrees in Z $< d^{2\varepsilon}$

EX $\Rightarrow |Z| \approx N/2$

\Rightarrow # of $I \cap Z$'s $\lesssim 2^{N/4}$

TRIV: G bipartite on $A \cup B \Rightarrow$

$$\text{m.i.s.}(G) \leq \min\{2^{|A|}, 2^{|B|}\}$$

\square

Non-example (?): # of k -SAT fns*
↙ ↘ more or less

$\mathcal{C} \subseteq \{ \text{ } \underline{k\text{-clauses}} \text{ from } x_1, \dots, x_n \}$
 e.g. $x_3 \bar{x}_{22} x_{91}$ ($k=3$)

with witnesses: $\forall C \in \mathcal{C} \exists w_C \in \{0,1\}^n$

$C \models w_C, C' \not\models w_C \forall C' \in \mathcal{C} \setminus \{C\}$

e.g. $x_1 x_2 \leftrightarrow (110 \dots 0)$
 $x_2 x_5 \leftrightarrow (010010 \dots 0)$
 $x_1 x_3 \leftrightarrow (1010 \dots 0) \dots$

$\rightarrow \left[\begin{array}{l} 2^{\binom{n}{k}} \text{ } \mathcal{C}'\text{'s} \\ \text{better: } 2^{n + \binom{n}{k}} \text{ } \mathcal{C}'\text{'s} \end{array} \right.$ ↗ choose literals

* Bollobás-Brightwell-Leader,
 P. Allen, Ilinca-K.

Conj 1 $\log \sim \binom{n}{k}$

OPEN $k > 3$

or

Conj 2 $\# \sim 2^{n + \binom{n}{k}}$

TRUE $k=2,3$

⇒ simple ("sampling") proof of

Conj 1 for $k=2$?