

# POLYNOMIAL-TIME PERFECT MATCHINGS IN DENSE HYPERGRAPHS

PETER KEEVASH, FIACHRA KNOX AND RICHARD MYCROFT

In this talk we consider the decision problem for the existence of a perfect matching in a  $k$ -uniform hypergraph (or  $k$ -graph)  $H$  on  $n$  vertices. Since for  $k \geq 3$  this problem was one of Karp's 21 NP-complete problems [1], it is natural to seek conditions on  $H$  which render it tractable. For any  $A \subseteq V(H)$ , the *degree*  $d(A) = d_H(A)$  of  $A$  is the number of edges of  $H$  containing  $A$ . The *minimum  $(k-1)$ -degree*  $\delta_{k-1}(H)$  of  $H$  is the minimum of  $d(A)$  over all subsets  $A$  of  $V(H)$  of size  $k-1$ .

Let  $\mathbf{PM}(k, \delta)$  be the decision problem of determining whether a  $k$ -graph  $H$  with  $\delta_{k-1}(H) \geq \delta n$  contains a perfect matching. Szymańska [3] proved that for  $\delta < 1/k$  the problem  $\mathbf{PM}(k, 0)$  admits a polynomial-time reduction to  $\mathbf{PM}(k, \delta)$  and hence  $\mathbf{PM}(k, \delta)$  is also NP-complete. We describe an algorithm which shows that the opposite is true for  $\delta > 1/k$ :

**Theorem 1.** *Fix  $k \geq 3$  and  $\gamma > 0$ . Then there is an algorithm with running time  $O(n^{3k^2-7k+1})$ , which given any  $k$ -graph  $H$  on  $n$  vertices with  $\delta_{k-1}(H) \geq (1/k + \gamma)n$ , finds either a perfect matching or a certificate that no perfect matching exists.*

Previously, Karpiński, Ruciński and Szymańska [2] showed that there exists  $\varepsilon > 0$  such that  $\mathbf{PM}(k, 1/2 - \varepsilon)$  is in P.

To prove Theorem 1 we establish a strong stability result which states that if  $H$  is a  $k$ -graph on  $n$  vertices, and  $\delta_{k-1}(H) \geq n/k + o(n)$ , then  $H$  either contains a perfect matching or is close to one of a family of lattice-based constructions termed ‘divisibility barriers’. While the precise statement of this result for general  $k$  requires significant preliminaries, which we cover in the talk, the special case  $k = 3$  may be stated as follows:

**Theorem 2.** *For any  $\gamma > 0$  there exists  $n_0 = n_0(\gamma)$  such that the following statement holds. Let  $H$  be a 3-graph on  $n \geq n_0$  vertices, such that 3 divides  $n$  and  $\delta_2(H) \geq (1/3 + \gamma)n$ , and suppose that  $H$  does not contain a perfect matching. Then there is a subset  $A \subseteq V(H)$  such that  $|A|$  is odd but every edge of  $H$  intersects  $A$  in an even number of vertices.*

## REFERENCES

- [1] R. M. Karp, Reducibility among combinatorial problems, *Complexity of Computer Computations* (1972), 85–103.
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- [3] E. Szymańska, The complexity of almost perfect matchings and other packing problems in uniform hypergraphs with high codegree, *European Journal of Combinatorics* **34** (2013), 632–646.

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