

Does branching explain flow of time or is it the other way around?

Petr Švarný

Logic and Relativity, 2012



Outline

- 1 Setup
 - Problem
 - Flow of time and branching models
- 2 The formalized interactions
 - FoT gives BCont a valuation
 - Example
 - Branches make FoT natural



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B-MODELS & FoT

The venture point

- Generalized flow of time can be formalized in Branching continuations and it adds the possibility of semantics to it.
- Generalized flow of time uses the term ontological definiteness as a basic notion.



B-MODELS & FoT

The interaction of branching models and flow of time

- (Does the framework of **Branching Continuations** entail the notion of a **generalized flow of time**?
- - OR -
- Does a generalized flow of time lead to a branching structure?)
- - AND -
- Does it help to solve the problem of now?



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GENERALIZED FLOW OF TIME

A generalized flow of time must have:

- worldlines with linear order of now-points
- ontological definiteness of past and present (!)
- a set of now-points on worldlines respecting the ontological definiteness

D. Dieks: Special relativity and the flow of time, 1988



BRANCHING MODELS

Main ideas of Branching Continuations

- $(W, \leq)$
- Continuations, I-events
- $e_C, e/A$
- Branching Time+*Instants*-like models

T. Placek: Possibilities Without Possible Worlds/Histories, 2009



BCONT MODEL

- 1 $\mathcal{W}$ is a non-empty partially ordered set;
- 2 the ordering $\leq$ is dense on W ;
- 3 W has no maximal elements;
- 4 every lower bounded chain $C \subseteq W$ has an infimum;
- 5 if a chain $C \subseteq W$ is upper bounded and $C \leq b$ then there is a unique minimum in $\{e \in W \mid C \leq e \wedge e \leq b\}$;
- 6 for every $x, y, e \in W$, if $e \not\leq x$ and $e \not\leq y$ then x and y are snake-linked in the subset $W_{\not\leq e} := \{e' \in W \mid e \not\leq e'\}$ of W ;
- 7 if $x, y \in W$ and $W_{\leq xy} := \{e \in W \mid e \leq x \wedge e \leq y\} \neq \emptyset$ then $W_{\leq xy}$ has a maximal element;
- 8 for every $x_1, x_2 \in W$, if $\forall c : c \in CE \rightarrow c \not\leq x_i$ then x_1, x_2 are snake-linked in the subset $W_{\not\leq CE} := \{e \in W \mid \forall c \in CE e \not\leq c\}$ of W .



BT+I MODEL

A model $\langle W, \leq, S \rangle$ is said to be *(BT+Instants)-like* if it satisfies the following conditions:

- downward directedness,
- no backward forks,
- $\forall e, e' \in W$: if e, e' are incomparable by $\leq$, then there are $H_1, H_2 \in \Pi_m$ such that $H_1 \neq H_2$, $e \in H_1$ and $e' \in H_2$, where m is a maximal element of $W_{\leq ee'} = \{y | y \leq e \wedge y \leq e'\}$;



GFOT IN BCONT

• Formalization of general FoT:

- worldline as chains of events
- settings of now-points
- ontological definiteness via valuation

P. Švarný: Flow of Time in BST/BCont Models and Related Semantical Observations, LOGICA 2012



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PFF

For given $e_C, e/A$ and the model $\mathfrak{M} = \langle \mathfrak{G}, \mathcal{I} \rangle$, then:

- if ψ is $P_x\varphi$ for $x > 0$: $\mathfrak{M}, e_C, e/A, X_{Wl(e_C),A} \Vdash \psi$ iff there is $e' \in \bigcup X_{Wl(e_C),A}$ such that $e' \cup A \in \text{I-events}$ and $\text{int}(e', e, Wl(e), x)$ and $\mathfrak{M}, e_C, e'/A \Vdash \varphi$;
- if ψ is $\text{Sett} : \varphi : \mathfrak{M}, e_C, e/A, X_{Wl(e_C),A} \Vdash \psi$ iff for every evaluation point e/A' from fan $\mathcal{F}_{e/A} : \mathfrak{M}, e_C, e/A', X_{Wl(e_C),A} \Vdash \varphi$;
- if ψ is $\text{Now} : \varphi : \mathfrak{M}, e_C, e/A, X_{Wl(e_C),A} \Vdash \psi$ iff there is $e' \in X_{e_C,A}$ such that $e' \cup A \in \text{I-events}$ and $\mathfrak{M}, e_C, e'/A, X_{Wl(e_C),A} \Vdash \varphi$.

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SETTLEDNESS OF THE FUTURE

Formulation of the problem

Is there a difference between necessary future events and future events?

$$\mathfrak{M}, e_C, e_C/A, X_{Wl(e_C),A} \models F_1\psi$$

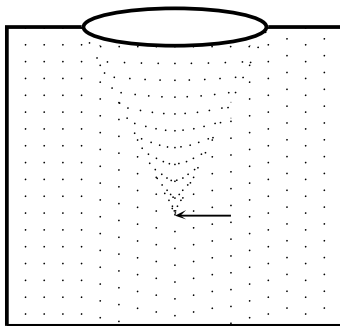
vs

$$\mathfrak{M}, e_C, e_C/A, X_{Wl(e_C),A} \models \text{Sett} : F_1\psi$$



SETTLEDNESS OF THE FUTURE I

Choice event in BCont+FoT

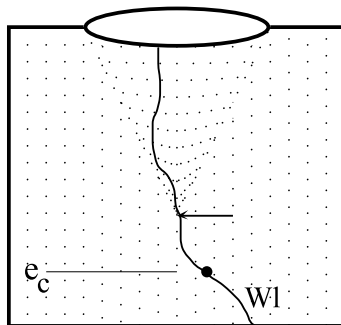


A 2D BCont model with a CE.



SETTLEDNESS OF THE FUTURE II

Worldlines in BCont+gFoT

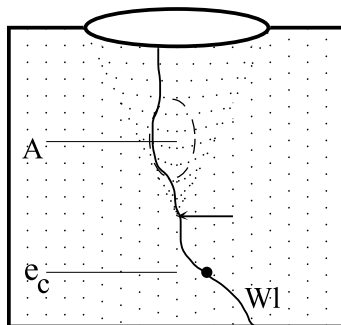


Model with a worldline and moment of use.



SETTLEDNESS OF THE FUTURE III

Continuations in BCont+gFoT

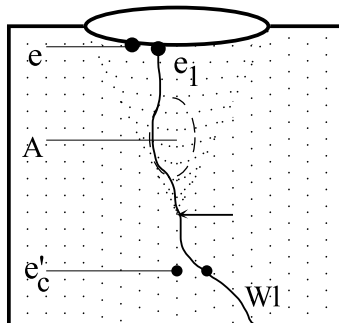


Adding the continuation A.



SETTLEDNESS OF THE FUTURE IV

Continuations in BCont+gFoT

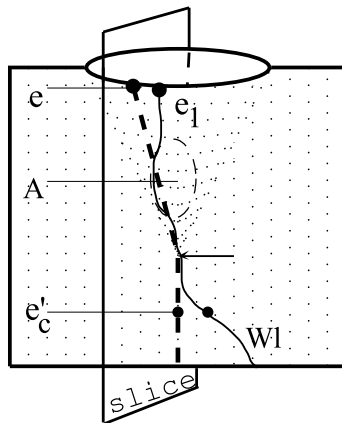


Adding the reference point e_1 and the point e that belongs to e_1 's setting of now-points and where ψ holds.



SETTLEDNESS OF THE FUTURE V

Continuations in BCont+gFoT

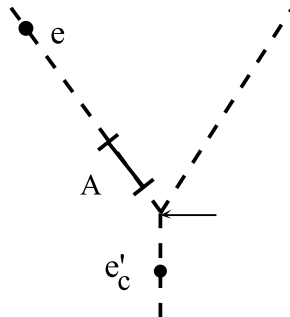
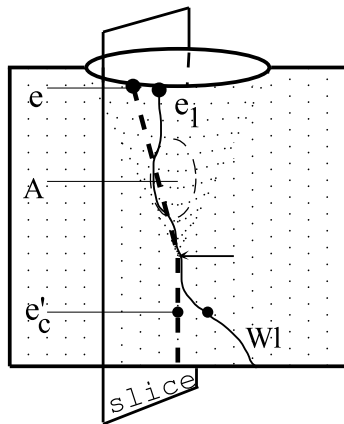


Preparing a slice going through point e .



SETTLEDNESS OF THE FUTURE VI

Continuations in BCont+gFoT



The final figure for *Sett* : $F_1\psi \not\equiv F_1\psi$.



GFOT IN BCONT

GFoT with respect to BCont has the following properties:

- is a definition away from basic BCont structures;
- allows for a valuation on BCont structures.



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GFOT DEMANDS

Let $gFoT$ a structure based on gFoT demands:

- structure is not necessarily BCont model (already ax. 2, nor BST).
- does not have to be any branching structure.
- difference of past and future given by now-points.



TRIVIAL BCONT

Let $\mathfrak{M} = \langle \mathfrak{G}, \mathcal{I} \rangle$ BCont model without CE:

- GFoT demands are met.
- Future is Peircean.
- Past/future has no significant difference.



FULL BCONT

Let $\mathfrak{M} = \langle \mathcal{G}, \mathcal{I} \rangle$ BCont model without CE:

- GFoT demands are met.
- Future is not Peircean.
- Past/Future are structurally different.



BCONT FOR GFOT

- GFoT does not need in any way BCont.
- BCont without choice events also fulfils gFoT's demands, but ontologically definite future.
- BCont with branches allows for a structural difference F/P and non-Peircean future.



SUMMARY

- GFoT is part of BCont structures and allows to construct valuation.
- GFoT does not need a branching structure. However, branching does make it non-trivial.
- Outlook:
 - analysis of temporal Copernican principle (Now).



Thank you for your attention.

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